# High-dimensional graphical model selection: Practical and information-theoretic limits

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#### Introduction

- classical asymptotic theory of statistical inference:
  - number of observations  $n \to +\infty$
  - model dimension p stays fixed
- not suitable for many modern applications:
  - { images, signals, systems, networks } frequently large  $(p\approx 10^3-10^8)...$
  - function/surface estimation: enforces limit  $p \to +\infty$
  - interesting consequences: might have  $p = \Theta(n)$  or even  $p \gg n$
- curse of dimensionality: frequently impossible to obtain consistent procedures unless  $p/n \to 0$
- can be saved by a lower *effective dimensionality*, due to some form of complexity constraint:
  - sparse vectors
  - {sparse, structured, low-rank}-matrices
  - structured regression functions
  - graphical models (Markov random fields)

#### What are graphical models?

• Markov random field: random vector  $(X_1, \ldots, X_p)$  with distribution factoring according to a graph G = (V, E):



• Hammersley-Clifford Theorem:  $(X_1, \ldots, X_p)$  being Markov w.r.t G implies factorization:

$$\mathbb{P}(x_1,\ldots,x_p) \propto \exp\left\{\theta_A(x_A) + \theta_B(x_B) + \theta_C(x_C) + \theta_D(x_D)\right\}.$$

• studied/used in various fields: spatial statistics, language modeling, computational biology, computer vision, statistical physics ....

#### Graphical model selection

• let G = (V, E) be an undirected graph on p = |V| vertices



• pairwise Markov random field: family of prob. distributions

$$\mathbb{P}(x_1,\ldots,x_p;\theta) = \frac{1}{Z(\theta)} \exp \Big\{ \sum_{(s,t)\in E} \langle \theta_{st}, \phi_{st}(x_s,x_t) \rangle \Big\}.$$

- given n independent and identically distributed (i.i.d.) samples of  $X = (X_1, \ldots, X_p)$ , identify the underlying graph structure
- complexity constraint: restrict to subset  $\mathcal{G}_{d,p}$  of graphs with maximum degree d



# Some issues in high-dimensional inference

Consider some fixed loss function, and a fixed level  $\delta$  of error.

#### Limitations of tractable algorithms:

Given particular (polynomial-time) algorithms

- for what sample sizes n do they succeed/fail to achieve error  $\delta$ ?
- given a collection of methods, when does more computation reduce minimum # samples needed?

#### Information-theoretic limitations:

Data collection as communication from nature  $\longrightarrow$  statistician:

- what are fundamental limitations of problem (Shannon capacity)?
- when are known (polynomial-time) methods optimal?
- when are there gaps between poly.-time methods and optimal methods?

# Previous/on-going work on graph selection

• exact solution for trees

(Chow & Liu, 1967)

local testing-based approaches
Buhlmann, 2008)

(e.g., Spirtes et al, 2000; Kalisch &

- methods for Gaussian MRFs
  - $\ell_1$ -regularized neighborhood regression for Gaussian MRFs (e.g., Meinshausen & Buhlmann, 2005; Wainwright, 2006, Zhao, 2006)
  - $\ell_1$ -regularized log-determinant (e.g., Yuan & Lin, 2006; d'Asprémont et al., 2007; Friedman, 2008; Ravikumar et al., 2008)
- methods for discrete MRFs
  - neighborhood-based search method
  - $-\ell_1$ -regularized logistic regression
- information-theoretic approaches:
  - pseudolikelihood and BIC criterion
  - information-theoretic limitations

(Bresler, Mossel & Sly, 2008) (Ravikumar et al., 2006, 2008)

(Csiszar & Talata, 2006)

(Santhanam & Wainwright, 2008)

# Markov property and neighborhood structure

• Markov properties encode neighborhood structure:



• basis of pseudolikelihood method



#### Practical method via neighborhood regression

**Observation:** Recovering graph G equivalent to recovering neighborhood set N(r) for all  $r \in V$ .

**Method:** Given *n* i.i.d. samples  $\{X^{(1)}, \ldots, X^{(n)}\}$ , perform logistic regression of each node  $X_r$  on  $X_{\backslash r} := \{X_r, t \neq r\}$  to estimate neighborhood structure  $\widehat{N}(r)$ .

1. For each node  $r \in V$ , perform  $\ell_1$  regularized logistic regression of  $X_r$  on the remaining variables  $X_{\backslash r}$ :

$$\widehat{\theta}[r] := \arg \min_{\theta \in \mathbb{R}^{p-1}} \left\{ \underbrace{\frac{1}{n} \sum_{i=1}^{n} f(\theta; X_{\backslash r}^{(i)})}_{\text{logistic likelihood}} + \rho_n \underbrace{\|\theta\|_1}_{\text{regularization}} \right\}$$

- 2. Estimate the local neighborhood  $\widehat{N}(r)$  as the support (non-negative entries) of the regression vector  $\widehat{\theta}[r]$ .
- 3. Combine the neighborhood estimates in a consistent manner (AND, or OR rule).

# **High-dimensional analysis**

- classical analysis: dimension p fixed, sample size  $n \to +\infty$
- high-dimensional analysis: allow both dimension p, sample size n, and maximum degree d to increase at arbitrary rates



- take n i.i.d. samples from MRF defined by  $G_{p,d}$
- study probability of success as a function of three parameters:  $Success(n, p, d) = \mathbb{P}[Method recovers graph G_{p,d} from n samples]$
- theory is non-asymptotic: explicit probabilities for finite (n, p, d)





#### Sufficient conditions for consistent model selection

- graph sequences  $G_{p,d} = (V, E)$  with p vertices, and maximum degree d.
- drawn n i.i.d, samples, and analyze prob. success indexed by (n, p, d)

**Theorem:** For a rescaled sample size (RavWaiLaf06, RavWaiLaf08)  $T_{\rm LR}(n, p, d) := \frac{n}{d^3 \log n} > T_{\rm crit}^*$ and regularization parameter  $\rho_n \ge c_1 \tau \sqrt{\frac{\log p}{n}}$ , then with probability greater than  $1 - 2 \exp\left(-c_2(\tau - 2)\log p\right) \rightarrow 1$ : (a) For each node  $r \in V$ , the  $\ell_1$ -regularized logistic convex program has a unique solution. (Non-trivial since  $p \gg n \Longrightarrow$  not strictly convex). (b) The estimated sign neighborhood  $\widehat{N}_{\pm}(r)$  correctly excludes all edges *not* in the true neighborhood.

(c) For  $\theta_{min} \geq c_3 \tau \sqrt{\frac{d^2 \log p}{n}}$ , the method selects the correct signed neighborhood.

# Some challenges in distinguishing graphs A B Hidden interactions Guilt by association Conditions on Fisher information matrix $Q^* = \mathbb{E}[\nabla^2 f(\theta^*; X)]$

- A1. Bounded eigenspectra:  $\lambda(Q_{SS}^*) \in [C_{min}, C_{max}].$
- A2. Mutual incoherence There exists an  $\nu \in (0, 1]$  such that

$$|||Q_{S^c S}^*(Q_{SS}^*)^{-1}|||_{\infty,\infty} \le 1-\nu.$$

where  $|||A|||_{\infty,\infty} := \max_i \sum_j |A_{ij}|.$ 

#### **Proof sketch: Primal-dual certificate**

- construct candidate primal-dual pair  $(\hat{\theta}, \hat{z}) \in \mathbb{R}^{p-1} \times \mathbb{R}^{p-1}$ .
- proof technique—-not a practical algorithm!

(A) For a fixed node r with S = N(r), we solve the restricted program  $\widehat{\theta} = \arg \min_{\theta \in \mathbb{R}^{p-1}, \theta_{S^c} = 0} \left\{ \frac{1}{n} \sum_{i=1}^n f(\theta; X_{\backslash r}^{(i)}) + \rho_n \|\theta\|_1 \right\},$ thereby obtaining candidate solution  $\widehat{\theta} = (\widehat{\theta}_S, \overrightarrow{0}_{S^c}).$ (B) We choose  $\widehat{z}_S \in \mathbb{R}^{|S|}$  as an element of the subdifferential  $\partial \|\widehat{\theta}_S\|_1$ . (C) Using optimality conditions from original convex program, solve for  $\widehat{z}_{S^c}$  and check whether or not strict dual feasibility  $|\widehat{z}_j| < 1$  for all  $j \in S^c$  holds.

**Lemma:** Full convex program recovers neighborhood  $\iff$  primal-dual witness succeeds.

#### Information-theoretic limits on graph selection

• thus far: have exhibited a a particular polynomial-time method can recover structure if

$$n > \Omega(d^3 \log(p - d))$$

- but....is this a "good" result?
- are there polynomial-time methods that can do better?
- information theory can answer the question: is there an exponential-time method that can do better?

(Santhanam & Wainwright, 2008)

# Graph selection as channel coding

- graphical model selection is an *unorthodox* channel coding problem:
- nature sends  $G \in \mathcal{G}_{d,p} := \{ \text{ graphs on } p \text{ vertices, max. degree } d \}$



- decoding problem: use observations  $\{X^{(1)}, \ldots, X^{(n)}\}$  to correctly distinguish the "codeword"
- channel capacity for graph decoding: balance between
  - log number of models:  $\log |M(p,d)| = \Theta \left( pd \log \frac{p}{d} \right).$
  - relative distinguishability of different models

#### Necessary conditions for graph recovery

- take Ising models  $\mathbb{P}_{\theta(G)}$  from  $\mathcal{G}_{d,p}(\lambda,\omega)$ :
  - graphs with p nodes and max. degree d
  - parameters  $|\theta_{st}| \ge \lambda$  for all edges (s, t)
  - maximum neighborhood weight  $\omega = \max_{s \in V} \sum_{t \in N(s)} |\theta_{st}|.$
- take n i.i.d. observations, and study probability of success in terms of (n, p, d)

**Theorem: Necessary conditions:** For sample size n

$$n \leq \max\left\{\frac{\log p}{2\lambda \tanh(\lambda)}, \frac{\exp(\omega/2)\lambda}{16\sinh(\lambda)} \ d\log(pd), \ \frac{d}{8}\log\frac{p}{8d}, \right\},$$

then the probability of error of any algorithm over  $\mathcal{G}_{d,p}(\lambda,\omega)$  is at least 1/2.

(Santhanam & W., 2008)

#### Some consequences

- note neighborhood weight  $\omega = \max_{s \in V} \sum_{t \in N(s)} |\theta_{st}|$  is at least  $d\lambda$
- hence, need at least

$$n > \frac{\exp(\frac{d\lambda}{2})\lambda}{16\sinh(\lambda)} d\log(pd)$$

- if  $\lambda = \mathcal{O}(1/d)$ , then need at least  $n > \frac{\log p}{\lambda^2} = \Omega(d^2 \log p)$  samples
- $\ell_1$ -regularized log. regression (LR) order-optimal for constant degrees
- for d tending to infinity, gap between optimal methods and  $\ell_1$ 
  - any method requires  $n = \Omega(d^2 \log p)$  samples
  - LR method: guaranteed to work with  $n = \Omega(d^3 \log p)$  samples

# Geometric intuition underlying proofs $D_2$ ruth Error probability controlled by two competing quantities: Model type Log # modelsDistance scaling $c_2/\theta^2$ Near-by $\log p$ $\sinh(\theta)$ Intermediate $d\log p$ $\overline{\theta \exp(\theta d)}$ $pd\log\frac{p}{d}$ Far-away $c_2 p$

#### Summary and open questions

•  $\ell_1$ -regularized regression to select neighborhoods: succeeds with sample size

$$n > \left(\frac{c_1}{\theta_{min}^2} + c_2 d^3\right) \log p.$$

• any method (including those with exponential complexity) fails for

$$n < \left(\frac{c_3}{\theta_{min}^2} + c_4 d^2\right) \log p$$

- some extensions....
  - non-binary MRFs via block-structured regularization schemes
  - non-i.i.d. sampling models
  - other performance metrics (e.g,  $(1 \delta)$  edges correct)
- broader issue: optimal trade-offs between statistical/computational efficiency?

#### Some papers

- Ravikumar, P., Wainwright, M. J. and Lafferty, J. (2008). High-dimensional Ising model selection using l<sub>1</sub>-regularized logistic regression. Appeared at NIPS Conference (2006); To appear in Annals of Statistics.
- Santhanam, P. and Wainwright, M. J. (2008). Information-theoretic limitations of high-dimensional graphical model selection. Presented at Int. Symposium on Information Theory.
- Wainwright, M. J. (2006). Sharp thresholds for noisy and high-dimensional recovery of sparsity using l<sub>1</sub>-constrained quadratic programming. To appear in *IEEE Trans. on Information Theory.*
- Wainwright, M. J. (2007). Information-theoretic limits on sparsity recovery in the high-dimensional and noisy setting. UC Berkeley, Department of Statistics, Technical Report, January 2007.