# Singularity Degree of PSD Matrix Completion 

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## Positive Semidefinite Matrix Completion

PSD completion problem ( $G, c$ )
Given $G=(V, E)$ with $V=\{1, \ldots, n\}$ and edge weight $c: E \rightarrow[-1,1]$,

$$
\begin{array}{ccc}
\text { find } & X \in \mathcal{S}^{n} & \\
\text { s.t. } & X[i, j]=c(i j) & (i j \in E) \\
& X[i, i]=1 & (i \in V) \\
& X \succeq 0 &
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$$
\begin{array}{cc}
\min & \langle\Omega, C\rangle \\
\text { s.t. } & \Omega \in S_{+}(G)
\end{array}
$$

where

$$
\begin{aligned}
C[i, j] & = \begin{cases}c(i j) & (i j \in E) \\
1 & (i=j) \\
0 & (\text { otherwise })\end{cases} \\
S(G) & :=\left\{A \in \mathcal{S}^{n}: A[i, j]=0 \forall i j \notin V \cup E\right\} \\
S_{+}(G) & :=\{A \in S(G): A \succeq 0\} \\
\bar{S}_{+}(G) & :=\left\{A \in S_{+}(G): A[i, j] \neq 0 \forall i j \in E\right\}
\end{aligned}
$$

## Geometric View

- Given a completion problem ( $G, c$ ),
- PSD completion $X=P P^{\top}$ with rank $d$ $\Leftrightarrow$ spherical embedding $p: V \rightarrow \mathbb{S}^{d-1}$ realizing $c$, i.e.,

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p_{i} \cdot p_{j}=c(i j) \quad \forall i j \in E
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\begin{align*}
\langle C, \Omega\rangle=0 & \Leftrightarrow\langle X, \Omega\rangle=0 \\
& \Leftrightarrow \Omega P=0 \\
& \Leftrightarrow \Omega[i, i] p(i)+\sum_{j \sim i} \Omega[i, j] p(j)=0 \quad(\forall i \in V) \tag{1}
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$$

- $\Omega$ is called a stress (matrix) of ( $G, p$ ) if $\Omega$ satisfies (1)
- Given $(G, p), \Omega \in S(G)$ is dual opt iff $\Omega$ is a PSD stress of $(G, p)$.


## SDP Duality

For any primal and dual optimal pair $(X, \Omega)$,

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\langle X, \Omega\rangle=0 \Rightarrow \operatorname{rank} X+\operatorname{rank} \Omega \leq n .
$$

- high rank dual opt $\Rightarrow$ low rank completion


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## Rank maximality certificate

- A completion $X$ for $(G, c)$ attains the maximum rank if $\exists$ dual opt with rank $n-\operatorname{rank} X$.


## Parameter $\nu$ and Unique Completability

## Theorem (Connelly82, Laurent-Varvitsiotis14)

- A completion $X$ for $(G, c)$ is unique if $\exists$ dual opt $\Omega$ with $\operatorname{rank} \Omega=n-\operatorname{rank} X$ and the SAP, i.e.,

$$
\nexists X \in \mathcal{S}^{n} \backslash\{0\} \text { with } \Omega X=0 \text { and } X[i, j]=0 \text { for } i j \in W E
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- $(G, p)$ is universally rigid in $\mathbb{S}^{d-1}$ if $(G, p)$ admits a PSD stress $\Omega$ with $\operatorname{rank} \Omega=n-d$ and the SAP.


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Colin de Verdière Parameter $\nu$

$$
\nu(G):=\max \left\{\operatorname{corank} \Omega: \Omega \in \bar{S}_{+}(G) \text { has the SAP }\right\} .
$$

$\nu(G) \leq \max \left\{d: \exists\right.$ universally rigid $(G, p)$ in $\left.\mathbb{S}^{d-1}\right\}$

## Strict Complementarity and Singularity Degree

## Strict Complementarity

A primal and dual optimal pair $(X, \Omega)$ satisfies a strict complementarity condition if

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\operatorname{rank} X+\operatorname{rank} \Omega=n
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- For which problem the strict complementarity can be guaranteed?
- How far from the strict complementarity?


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## Proposition

The following are equivalent for a graph $G$ :
(1) $\operatorname{sd}(G)=1$;
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## Proposition

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(1) $\operatorname{sd}(G)=1$;
(2) The strict complementarity holds for any PSD completion problem with underlying graph $G$;
(0) The projection $\mathcal{E}(G)$ of the elliptope (the set of correlation matrices) onto $\mathbb{R}^{E}$ is exposed (Druvyatskiy-Pataki-Wolkowicz15).

## Facial Reduction (Borwein-Wolkowitcz81)

A sequence $\left\{\Omega_{1}, \ldots, \Omega_{k}\right\}$ in $\mathcal{S}^{n}$ is iterated PSD if $\Omega_{i}$ is positive semidefinite on $\mathcal{V}_{i-1}$, where $\mathcal{V}_{0}=\mathbb{R}^{n}$ and

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Theorem (Facial reduction)
For any feasible ( $G, c$ ), $\exists X$ and $\exists \Omega_{1}, \ldots, \Omega_{k} \in S(G)$ s.t.
(1) the sequence is iterated PSD
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## Definition (Sturm 2000)

For a completion problem $(G, c)$, the singularity degree $\operatorname{sd}(G, c)$ is the length of the shortest dual certificate sequence $\left\{\Omega_{1}, \ldots, \Omega_{k}\right\}$.

## Singularity Degree of Graphs

## Singularity degree of $G$

$$
\operatorname{sd}(G)=\max _{c} \operatorname{sd}(G, c)
$$

Question (Druvyatskiy-Pataki-Wolkowicz15) Characterize $G$ with $\operatorname{sd}(G)=1$ Question (So15) $\operatorname{sd}(G)=o(n)$ ?

## Main Results

Theorem (T16)
$\operatorname{sd}(G)=1$ iff $G$ is chordal.
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- If $G$ has neither $W_{n}(n \geq 5)$ nor a proper splitting of $W_{n}(n \geq 4)$ as an induced subgraph, then $\operatorname{sd}(G) \leq 2$.
- If $G$ has an induced subgraph which is a proper splitting of one of the above forbidden subgraphs, then $\operatorname{sd}(G)>2$.

If $\operatorname{tw}(G) \leq 2$, then $\operatorname{sd}(G) \leq 2$.

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If $\operatorname{tw}(G) \leq 2$, then $\operatorname{sd}(G) \leq 2$.

## Theorem (T16)

For each $n$ there is a graph $G$ with $n$ vertices and $\operatorname{tw}(G)=3$ whose singularity degree is $\left\lfloor\frac{n-1}{3}\right\rfloor$.

## Proof of the first theorem

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Lemma. $\operatorname{sd}(G) \geq \operatorname{sd}(H)$ for any induced subgraph $H$ of $G$.

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- Consider ( $G, p$ ):

- $(G, p)$ is universally rigid
- there is a unique stress $\Omega$ with $\operatorname{rank} \Omega=1<n-2$

Lemma. $\operatorname{sd}(G) \geq \operatorname{sd}(H)$ for any induced subgraph $H$ of $G$.

## Nondegenerate Singularity Degree

A completion problem $(G, c)$ is nondegenerate if $c(i j) \neq \pm 1$ for every $i j \in E(G)$.
Degenerate edges can easily be eliminated.

- Suppose $c(i j)=1$ for $i j \in E \ldots$
- Any solution $X$ of $(G, c)$ satisfies

$$
X[i, k]=X[j, k] \quad \text { for every } k
$$

- Equivalently, any embedding $p$ realizing $c$ satisfies $p(i)=p(j)$. The example in the last proof is degenerate...


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Lemma $\operatorname{sd}(G) \leq \operatorname{sd}^{*}(G)+1$.

Corollary (T16) $\operatorname{sd}(G) \leq 2$ if $G$ has no forbidden induced subgraph listed above.

## Example of Large Singularity Degree

## Theorem (T16)

For each $n$ there is a graph $G$ with $n$ vertices and $\operatorname{tw}(G)=3$ whose singularity degree is $\left\lceil\frac{n-1}{3}\right\rceil$.


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- Q. Characterize graphs $G$ with $\operatorname{sd}(G) \leq 2$.
- Q. Bound $\operatorname{sd}(G)$ by other graph parameters.

