

Graph Cores via Universal Completability

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Outline

- Universal completable frameworks
- Sufficient condition for UC
- Least eigenvalue frameworks
- Sufficient condition for showing a graph is a core

Universal completability

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ex: a C_{2k+1} -framework

$$p_x = \left(\cos\left(\frac{2\pi kx}{2k+1}\right), \sin\left(\frac{2\pi kx}{2k+1}\right) \right), \quad 0 \leq x \leq 2k,$$

Universal completability

Def: A framework $G(\mathbf{p})$ is called **UC** if $\forall q_1, \dots, q_n$

$$\langle p_i, p_j \rangle = \langle q_i, q_j \rangle \quad \text{for } [i = j \ \& \ ij \in E]$$

implies that

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→ the G -partial matrix corresponding to $G(\mathbf{p})$
has a **unique** PSD completion

Why is it interesting?

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→ find G -partial PSD matrix

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→ $\text{corank}(Z) = d$

→ $\text{span}(p_i p_j^T + p_j p_i^T : i, j \in V \cup E)$

How to use this?

Nullspace representations

Def: Let $M \in \mathcal{S}^n$ with $\text{corank}(M) = d$ and $\{u_i\}_{i=1}^d$ an orth. basis for $\text{Ker}(M)$. Set $P = [u_1, \dots, u_d]$ and let $\{p_i\}_{i=1}^n \subseteq \mathbb{R}^d$ be the rows of P . The map

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e.g. [Lovász, Schrijver'99, '06]



Constructing outerplanar embeddings

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→ Brouwer & Haemers (Euclidean reps.)

Least eigenvalue framework

Thm: A least eig. framework $G(p) \subseteq \mathbb{R}^m$ is UC iff

$$\forall R \in \mathcal{S}^m : p_i^\top R p_j = 0 \text{ for } i = j \text{ and } (i, j) \in E \implies R = 0.$$

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Pf: \Leftarrow Use the sufficient condition for UC with

$$Z = A - \lambda_{\min}(A)I$$

\implies Scale so $\lambda_{\min}(R) = -1$ and note that

$$P(I + R)P^T \neq PP^T$$

Graph cores

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→ Complete graphs, odd cycles are cores

→ Bipartite graphs are not cores

Cores

Def: An endomorphism f is **locally injective** if

$$d(i, j) = 2 \implies f(i) \neq f(j)$$

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Thm [Nešetřil]: Every locally injective endomorphism is an automorphism

Graph cores via UC

Thm: Consider a framework $G(\mathbf{p})$ with

→ $d(i, j) = 2 \implies p_i \neq p_j$

→ $\|p_i\|$ constant for $i \in [n]$

→ $\langle p_i, p_j \rangle$ constant for $ij \in E$

→ $G(\mathbf{p})$ is UC

Then G is a core.

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Pf: Show that any end. f is locally injective

Graph cores via UC

Pf: Say $\exists u, v$ such that $d(u, v) = 2$ and $f(u) = f(v)$

\rightarrow Set $p'_i = p_{f(i)}, \forall i$

$\rightarrow G(p)$ and $G(p')$ are equivalent

$\rightarrow \exists$ isometry U such that $p'_i = Up_i, \forall i$

So $\langle p_u, p_v \rangle = \langle p'_u, p'_v \rangle = \langle p_{f(u)}, p_{f(v)} \rangle$ which implies

$p_u = p_v$ contradiction

Graph cores via UC

Thm: If there exists a framework $G(\mathbf{p})$ such that:

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Use LEF

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→ $G(\mathbf{p})$ is UC

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What about these?

1-walk-regular graphs

Def: G is **1WR** if $\exists a_k, b_k \in \mathbb{N}$ such that:

$$\rightarrow A^k \circ I = a_k I$$

$$\rightarrow A^k \circ A = b_k A$$

for all $k \in \mathbb{N}$

Recall $(X \circ Y)_{ij} = X_{ij} Y_{ij}$

1-walk-regular graphs

Thm: Let G be 1WR and $G(p) \subseteq \mathbb{R}^m$ a LEF. Then:

$$\Rightarrow \|p_i\|^2 = \frac{m}{n} \text{ for all } i \in [n]$$

$$\Rightarrow \langle p_i, p_j \rangle = \frac{\tau m}{na_2} \text{ for all } ij \in E$$

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Pf: Let E_τ the projector on the least eigenspace

→ E_τ is a polynomial in A

→ $\exists a, b \in \mathbb{R}$ such that

$$E_\tau \circ I = aI \quad \text{and} \quad E_\tau \circ A = bA$$

1-walk-regular graphs

Pf: Let E_τ the projector on the least eigenspace

→ E_τ is a polynomial in A

→ $\exists a, b \in \mathbb{R}$ such that

$$E_\tau \circ I = aI \quad \text{and} \quad E_\tau \circ A = bA$$

→ $E_\tau = \text{Gram}(p_1, \dots, p_n)$

Putting everything together

Thm: Let G be 1WR and $G(p) \subseteq \mathbb{R}^m$ a LEF. Say that

$$\rightarrow d(i, j) = 2 \implies p_i \neq p_j$$

$$\rightarrow \text{span}(p_i p_j^T + p_j p_i^T : i, j \in V \cup E)$$

Then G is a core.

Applications

Kneser graphs

Def: The Kneser graph $K_{n:r}$ has:

Vertices: Subsets of $[n]$ of size r

Edges: $X \sim Y$ if $X \cap Y = \emptyset$

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Thm: [Hahn, Tardif'97]

For $n \geq 2r + 1$ the graph $K_{n:r}$ is a core

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Thm: [Hahn, Tardif'97]

For $n \geq 2r + 1$ the graph $K_{n:r}$ is a core

→ The proof uses EKR theorem

Kneser graphs

To show $K_{n:r}$ is a core using our sufficient condition:

→ Find LEF and show that

→ $p_i \neq p_j, \forall i, j$

→ $\text{span}(p_i p_j^T + p_j p_i^T : i, j \in V \cup E)$

References

- Universal completability, least eigenvalue frameworks and vector colorings, arXiv:1512.04972
- Vector colorings and graph homomorphisms, coming soon

Thank you!