Don't disturb my Flows: Consistent Migration of Flows in SDNs

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A Small Sample Network



Green wants to send as well



Congestion!



This would work



So lets go back



But Red is a bit Slow..



Congestion Again!





Appears in Practice



"some switches can '**straggle**,' taking substantially **more time** than average (e.g., 10-**100x**) to apply an update" Jin et al., SIGCOMM 2014

So lets go Back ...



First, Red switches



Then, Blue ...



And then, Green ...



Done

Introduced in SWAN (Hong et al., SIGCOMM 2013) Idea: Flows can be on the **old** or **new** route For all edges: $\sum_{\forall F} \max(\mathbf{old}, \mathbf{new}) \leq capacity$

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No ordering exists (2/3 + 2/3 > 1)



Approach of *SWAN*: use slack *x* (i.e., %)

Here x = 1/3Move slack $x \Rightarrow [1/x] - 1$ staged partial moves



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Update 2 of 2



No slack on flow edges?



Alternate routes?



Think: variable swapping of b & g1. $x \coloneqq b$, 2. $b \coloneqq g$, 3. $g \coloneqq x$



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SWAN: LP-approach with binary search 1 update? 2 updates? 4 updates? ...



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SWAN: LP-approach with binary search $\Theta(1/\varepsilon)$ updates \mathfrak{S}



Open problem: Can we decide in (polynomial) time?



To Slack or not to Slack?

Slack of x on all flow edges? [1/x] - 1 updates

To Slack or not to Slack?

What if not? Try to create slack

To Slack or not to Slack?

Combinatorial approach Augmenting paths

Move single commodities at a time



Where to increase flow?



Where to push back flow?



Resulting residual network



We found an augmenting path \Rightarrow create slack on e



High-level Algorithm Idea

No slack on flow edges? Find augmenting paths On both initial and desired state Success? Use SWAN method to migrate

Can't create slack on some flow edge?

Consistent migration impossible By contradiction (else augmenting paths would create slack)

Runtime: $O(Fm^3)$

(F being #commodities, m being #edges)

Are we done?

Consistent Migration = Lossless Migration?













Overview and Outlook

Consistent migration of flows Decidable in **polynomial** time

Flow migration "Halting Problem"

Lossless migration Fixed delays: NP-hard, Arbitrary delays: Polynomially decidable

Unsplittable/Integer flow migration: **NP-hard**



Migrate to new **demands**

Single-dest: **Polynomial** #updates, Multi-Dest: **Ongoing work**

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