UnivMon: Software-defined Monitoring with Universal Sketch

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#### **Typical Measurement Questions:**

- Who's sending a lot more traffic than 10min ago? (Change)
- Who's sending a lot from 10.0.1.0/16? (Heavy Hitter)
- Are you being DDoS-ed?

### Example: A Victim being DDoSed



# **Traditional: Packet Sampling**

Sample packets at random, aggregate into flows



Estimate: FSD, Entropy, Heavyhitters ...

#### Not good for fine-grained analysis Extensive literature on limitations for many tasks!

### **Application-Specific Sketches**



*Complexity*: Need per-metric implementation *Recent Example*: OpenSketch [NSDI'13] *Trend:* Many more applications appear!

#### A Generic Method



# Outline

- Motivation
- UnivMon System Design
- UnivMon Algorithm
- Evaluation

### Our Design: UnivMon



• **Late-binding** for applications: data plane is general-purpose.

• **One Sketch**: no need of memory allocation for multiple tasks

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- When g(x) = x<sup>0</sup>, G-sum is # of distinct items. [SuperSpreader]
- When g(x) = xlog(x), G-sum is the entropy norm. [Entropy]

### Basics: Streaming Algorithms

- 11513312465 .....(A stream S of length m with n unique items)
  - K-th Frequency Moments:

 $F_k(S) = \sum_{i=1}^n f_i^k$ 

 $F_0$  is n: the number of distinct items in S.

 $F_1$  is m: the length of S.

 $F_2$  is Gini's index of homogeneity.

When k>2, the space lower bound is  $\Omega(n^{1-2/k})$  [CKS'03] (For current applications, the case of k>2 is not that interesting)

• Heavy Hitters (Frequent Items):

g-heavy item i :  $g(f_i) > \alpha$  G-sum for some  $\alpha$ Count-min sketch [CM'04] is a popular L1-heavy hitter algorithm

### Universal Sketch Data Structure



#### **Estimating G-sum**



### Intuitions of Universal Sketch



Idea: Detect items that contribute most to G-sum

# Putting it together: UnivMon



## Apple Apple

![](_page_16_Figure_1.jpeg)

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# **Key Evaluation Questions**

- Feasibility of hardware deployment?
  ✓ Removed expensive data structures
  ✓ Implemented in P4
- UnivMon's accuracy
  - ✓ Compare to the up-to-date sketch algorithms
- UnivMon's stability
  ✓ Stabilities over different traces

## **Evaluation**

Comparison with custom sketches via OpenSketch

Single Route Peter Exalution (600 KB) ces)

![](_page_19_Figure_3.jpeg)

## **Evaluation**

Comparison with custom sketches via OpenSketch

Netwon Month And Network Network (600 K Backer sketch)

![](_page_20_Figure_3.jpeg)

### Conclusions

- Traditional packet sampling based approaches have limitations for fine-grained analysis.
- Custom Sketches (e.g. OpenSketch)
  - Need to know applications beforehand
  - High Implementation cost
  - More expensive on multi-tasks
- UnivMon: an efficient and general sketching based approach with late-binding on applications.

# **Future Directions**

- Multi-dimensional data
- Performance optimization for hardware
- Dynamically change monitoring scope
- New Theories on Universal Sketch:
  - Sliding Window Model [SODA'16]
  - Functions of One Variable [PODS'16]
  - Symmetric Norms [arXiv 1511.01111]

Thank you!