

Matrix Completion from Fewer Entries

Raghunandan Keshavan, Andrea Montanari and Sewoong Oh

Stanford University

March 30, 2009

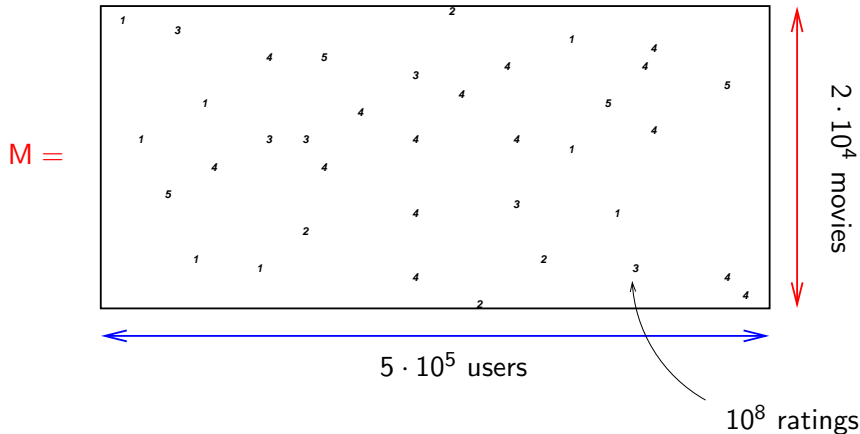
Outline

- 1 The problem, a look at the data, and some results (slides)
- 2 Proofs (blackboard)

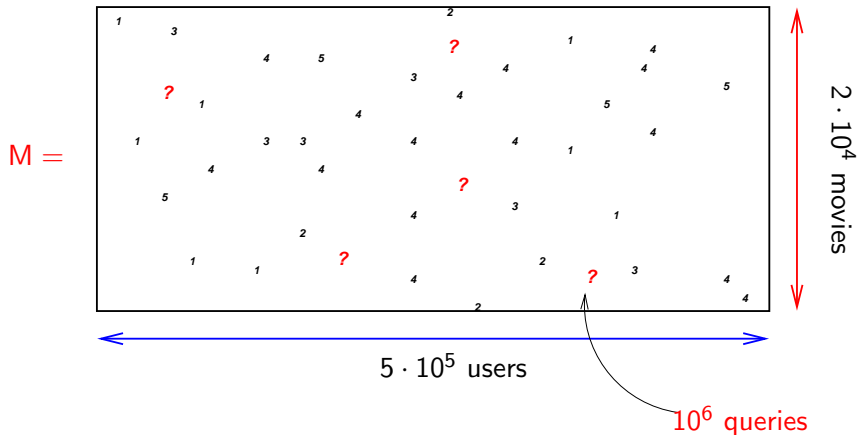
arXiv:0901.3150

The problem, a look at the data, and some results

Netflix dataset: A big (!) matrix



A big (!) matrix



You get a prize if...

RMSE < 0.8563 ; -)

Is this possible?

You get a prize if...

$\text{RMSE} < 0.8563$; -)

Is this possible?

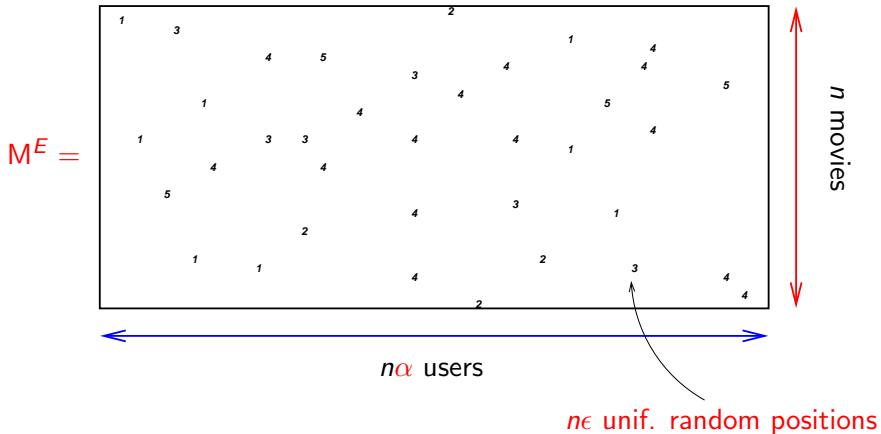
You get a prize if...

RMSE < 0.8563 ; -)

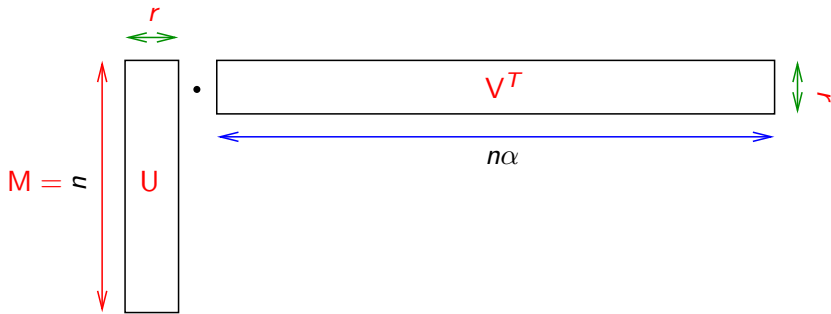
Is this possible?

A model: Incoherent low-rank matrices

The observations

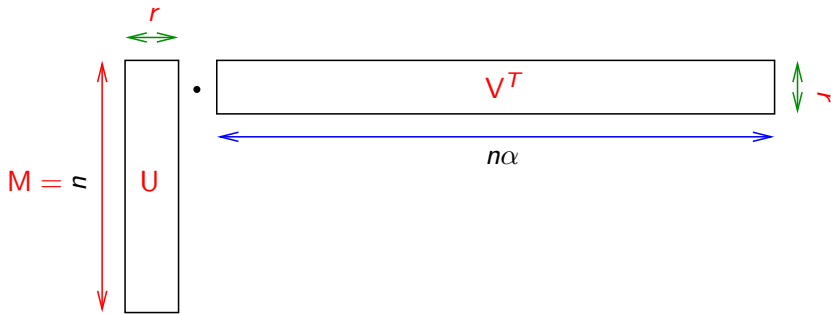


You need some structure!



$$r \ll n$$

You need some structure!



$$r \ll n$$

A1. Bounded entries

$$|M_{ia}| \leq M_{\max} = \mu_0 \sqrt{r}.$$

A2. Incoherence

$$\sum_{k=1}^r U_{ik}^2 \leq \mu_1 r, \quad \sum_{k=1}^r V_{ak}^2 \leq \mu_1 r.$$

[Candés, Recht 2008]

Metric (RMSE)

$$D(M, \hat{M}) \equiv \left\{ \frac{1}{n^2 M_{\max}^2} \sum_{i,a} |M_{ia} - \hat{M}_{ia}|^2 \right\}^{1/2}$$

Theorem (Candés, Recht, 2008)

If

$$\epsilon \geq C r n^{1/5} \log n$$

then whp

1. *M is unique given the observed entries.*
2. *M is the unique minimum of a SDP.*

cf. also [Recht, Fazel, Parrilo 2007]

Theorem (Candés, Recht, 2008)

If

$$\epsilon \geq C r n^{1/5} \log n$$

then whp

1. *M is unique given the observed entries.*
2. *M is the unique minimum of a SDP.*

cf. also [Recht, Fazel, Parrilo 2007]

Theorem (Candés, Recht, 2008)

If

$$\epsilon \geq C r n^{1/5} \log n$$

then whp

1. *M is unique given the observed entries.*
2. *M is the unique minimum of a SDP.*

cf. also [Recht, Fazel, Parrilo 2007]

Theorem (Candés, Recht, 2008)

If

$$\epsilon \geq C r n^{1/5} \log n$$

then whp

1. *M is unique given the observed entries.*
2. *M is the unique minimum of a SDP.*

cf. also [Recht, Fazel, Parrilo 2007]

Theorem (Candés, Recht, 2008)

If

$$\epsilon \geq C r n^{1/5} \log n$$

then whp

1. *M is unique given the observed entries.*
2. *M is the unique minimum of a SDP.*

cf. also [Recht, Fazel, Parrilo 2007]

Great, but...

1. $n^{1/5}$ observations for 1 bit of information?
2. RMSE = 0?
3. SDP = $O(n^{4\dots 6})$. Substitute $n = 10^5 \dots$

Great, but...

1. $n^{1/5}$ observations for 1 bit of information?
2. RMSE = 0?
3. SDP = $O(n^{4...6})$. Substitute $n = 10^5$...

Great, but...

1. $n^{1/5}$ observations for 1 bit of information?
2. RMSE = 0?
3. SDP = $O(n^{4...6})$. Substitute $n = 10^5$...

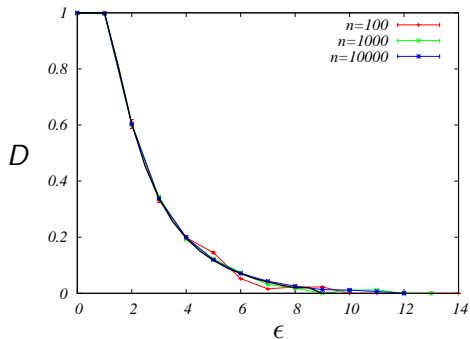
Great, but...

1. $n^{1/5}$ observations for 1 bit of information?
2. RMSE = 0?
3. SDP = $O(n^{4\dots 6})$. Substitute $n = 10^5 \dots$

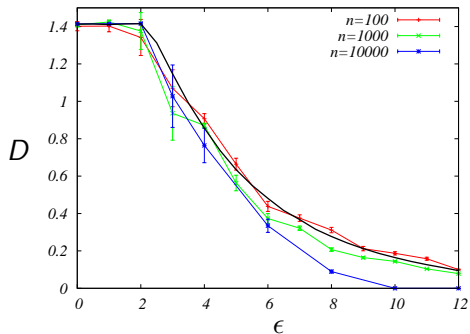
$O(n)$ entries are enough (practice)

A movie

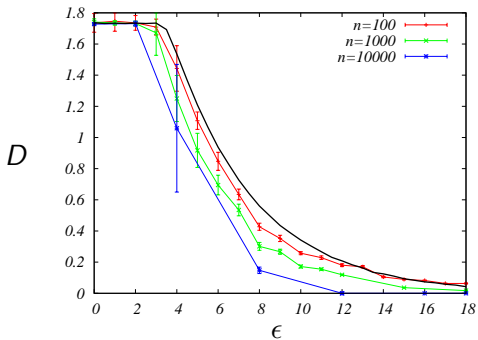
Rank = 1: Bayes optimal vs. Belief Propagation



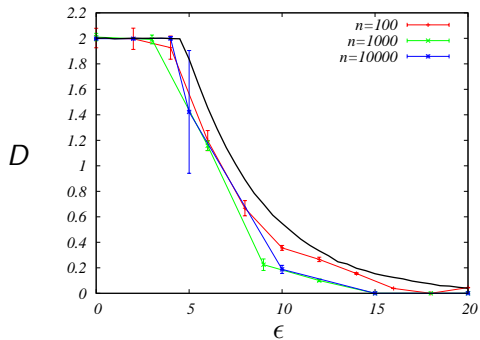
Rank = 2: Belief Propagation



Rank = 3: Belief Propagation



Rank = 4: Belief Propagation



$O(n)$ entries are enough (theory)

Naive spectral algorithm

$$M_{ia}^E = \begin{cases} M_{ia} & \text{if } (i, a) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

Projection

$$M^E = \sum_{i=1}^n \sigma_i x_i y_i^T, \quad \sigma_1 \geq \sigma_2 \geq \dots$$
$$\text{Tr}_r(M^E) = \frac{n\sqrt{\alpha}}{\epsilon} \sum_{i=1}^r \sigma_i x_i y_i^T.$$

Naive spectral algorithm

$$M_{ia}^E = \begin{cases} M_{ia} & \text{if } (i, a) \in E, \\ 0 & \text{otherwise.} \end{cases}$$

Projection

$$M^E = \sum_{i=1}^n \sigma_i x_i y_i^T, \quad \sigma_1 \geq \sigma_2 \geq \dots$$

$$\text{Tr}_r(M_E) = \frac{n\sqrt{\alpha}}{\epsilon} \sum_{i=1}^r \sigma_i x_i y_i^T.$$

If $\epsilon = O(1)$, 'spurious' singular values $\Omega(\sqrt{\log n / (\log \log n)})$.

Trimming

$$\tilde{M}_{ia}^E = \begin{cases} M_{ia}^E & \text{if } \deg(i) \leq 2 \mathbb{E} \deg(i), \quad \deg(a) \leq 2 \mathbb{E} \deg(a), \\ 0 & \text{otherwise.} \end{cases}$$

If $\epsilon = O(1)$, 'spurious' singular values $\Omega(\sqrt{\log n / (\log \log n)})$.

Trimming

$$\tilde{M}_{ia}^E = \begin{cases} M_{ia}^E & \text{if } \deg(i) \leq 2 \mathbb{E} \deg(i), \quad \deg(a) \leq 2 \mathbb{E} \deg(a), \\ 0 & \text{otherwise.} \end{cases}$$

Not-as-naive spectral algorithm

SPECTRAL MATRIX COMPLETION(matrix M^E)

- 1: Trim M^E , and let \tilde{M}^E be the output;
 - 2: Project \tilde{M}^E to $T_r(\tilde{M}^E)$;
 - 3: Clean residual errors by gradient descent in the factors.
-

Theorem (Keshavan, M, Oh, 2009)

Assume $r \leq n^{1/2}$ and bounded entries. Then

$$\frac{1}{nM_{\max}} \|\mathbf{M} - \mathbf{T}_r(\tilde{\mathbf{M}}^E)\|_F = \text{RMSE} \leq C\sqrt{r/\epsilon}.$$

with probability larger than $1 - \exp(-Bn)$.

Theorem (Keshavan, M, Oh, 2009)

Assume $r = O(1)$, bounded entries and incoherent factors, with $\Sigma_{\min}, \Sigma_{\max}$ uniformly bounded away from 0 and ∞ .

If $\epsilon \geq C' \log n$ then

SPECTRAL MATRIX COMPLETION returns, whp, the matrix \mathbf{M} .

Theorem (Keshavan, M, Oh, 2009)

Assume $r \leq n^{1/2}$ and bounded entries. Then

$$\frac{1}{nM_{\max}} \|\mathbf{M} - \mathbf{T}_r(\tilde{\mathbf{M}}^E)\|_F = \text{RMSE} \leq C\sqrt{r/\epsilon}.$$

with probability larger than $1 - \exp(-Bn)$.

Theorem (Keshavan, M, Oh, 2009)

Assume $r = O(1)$, bounded entries and incoherent factors, with $\Sigma_{\min}, \Sigma_{\max}$ uniformly bounded away from 0 and ∞ .

If $\epsilon \geq C' \log n$ then

SPECTRAL MATRIX COMPLETION returns, whp, the matrix \mathbf{M} .

A comparison

Theorem (Achlioptas, McSherry 2007)

Assume $\epsilon \geq (8 \log n)^4$ and bounded entries. Then

$$\frac{1}{nM_{\max}} \|\mathbf{M} - \text{Tr}(\tilde{\mathbf{M}}^E)\|_F = \text{RMSE} \leq 4\sqrt{r/\epsilon}.$$

with probability larger than $1 - \exp(-19(\log n)^4)$.

(For $n = 10^6$, $(8 \log n)^4 \approx 1.5 \cdot 10^8$)

Theorem (Achlioptas, McSherry 2007)

Assume $\epsilon \geq (8 \log n)^4$ and bounded entries. Then

$$\frac{1}{nM_{\max}} \|\mathbf{M} - \text{Tr}(\tilde{\mathbf{M}}^E)\|_F = \text{RMSE} \leq 4\sqrt{r/\epsilon}.$$

with probability larger than $1 - \exp(-19(\log n)^4)$.

(For $n = 10^6$, $(8 \log n)^4 \approx 1.5 \cdot 10^8$)

One more comparison

Theorem (Candés, Tao, March 8, 2009)

Assume bounded entries and *strongly incoherent factors*

If $\epsilon \geq C r (\log n)^6$ then

SEMIDEFINITE PROGRAMMING returns, whp, the matrix M .

A2'. Strong incoherence

$$\sum_{k=1}^r U_{ik}^2 \leq \mu_1 r,$$

$$\left| \sum_{k=1}^r U_{ik} U_{jk} \right| \leq \mu_1 \sqrt{r},$$

One more comparison

Theorem (Candés, Tao, March 8, 2009)

Assume bounded entries and *strongly incoherent factors*

If $\epsilon \geq C r (\log n)^6$ then

SEMIDEFINITE PROGRAMMING returns, whp, the matrix M .

A2'. Strong incoherence

$$\sum_{k=1}^r U_{ik}^2 \leq \mu_1 r,$$

$$\left| \sum_{k=1}^r U_{ik} U_{jk} \right| \leq \mu_1 \sqrt{r},$$

One more comparison

Theorem (Candés, Tao, March 8, 2009)

Assume bounded entries and *strongly incoherent factors*

If $\epsilon \geq C r (\log n)^6$ then

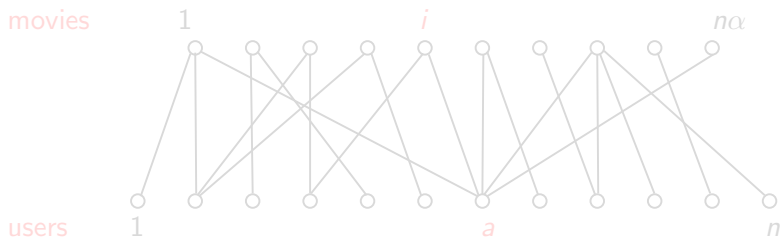
SEMIDEFINITE PROGRAMMING returns, whp, the matrix M .

A2'. Strong incoherence

$$\sum_{k=1}^r U_{ik}^2 \leq \mu_1 r,$$

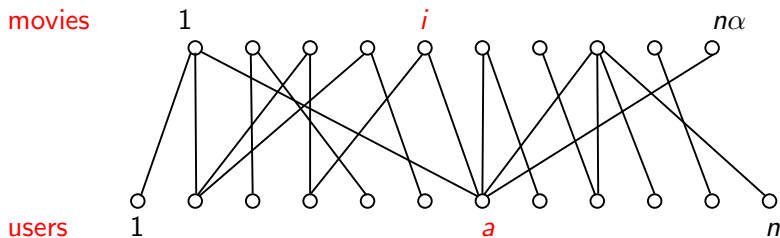
$$\left| \sum_{k=1}^r U_{ik} U_{jk} \right| \leq \mu_1 \sqrt{r},$$

Our approach: Graph theory



$(i, a) \in E \Leftrightarrow$ User a rated movie i .

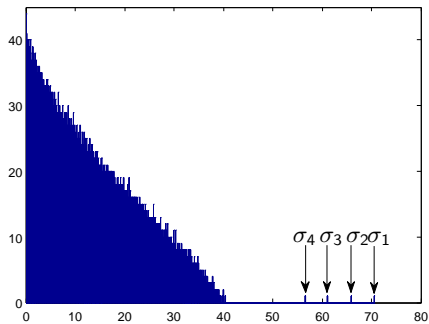
Our approach: Graph theory



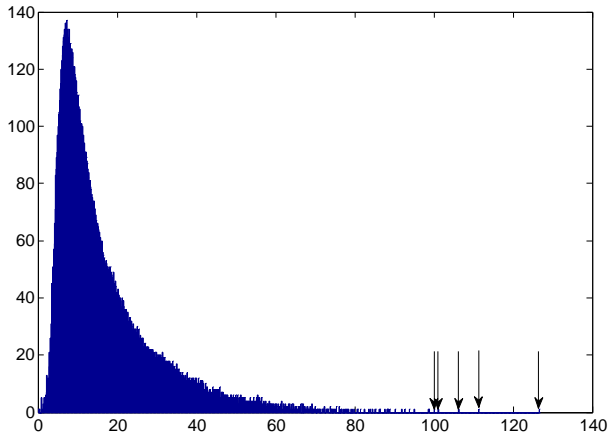
$(i, a) \in E \Leftrightarrow$ User a rated movie i .

Back to the data

Random $r = 4$, $n = 10000$, $\epsilon = 12.5$



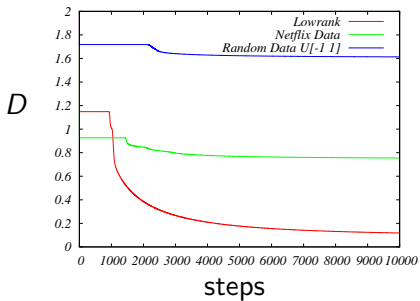
Netflix data (trimmed)



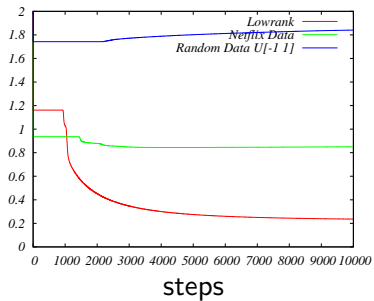
Is Netflix a random low-rank matrix?

Compare for coordinate descent (SimonFunk).

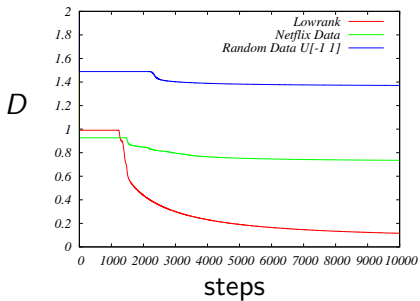
fit error



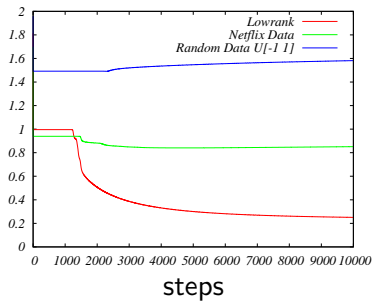
pred. error



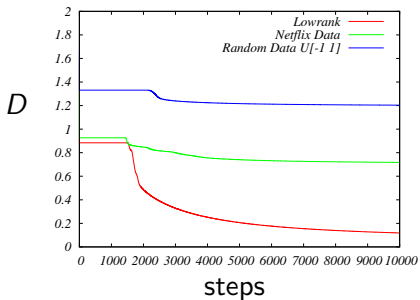
fit error



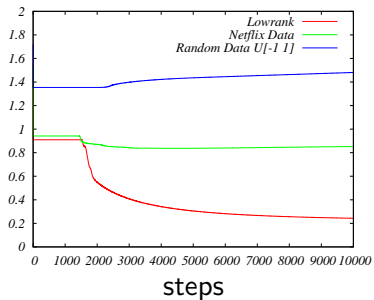
pred. error



fit error



pred. error



Proofs (blackboard)