

# Dynamic Pricing for Non-Perishable Products with Demand Learning

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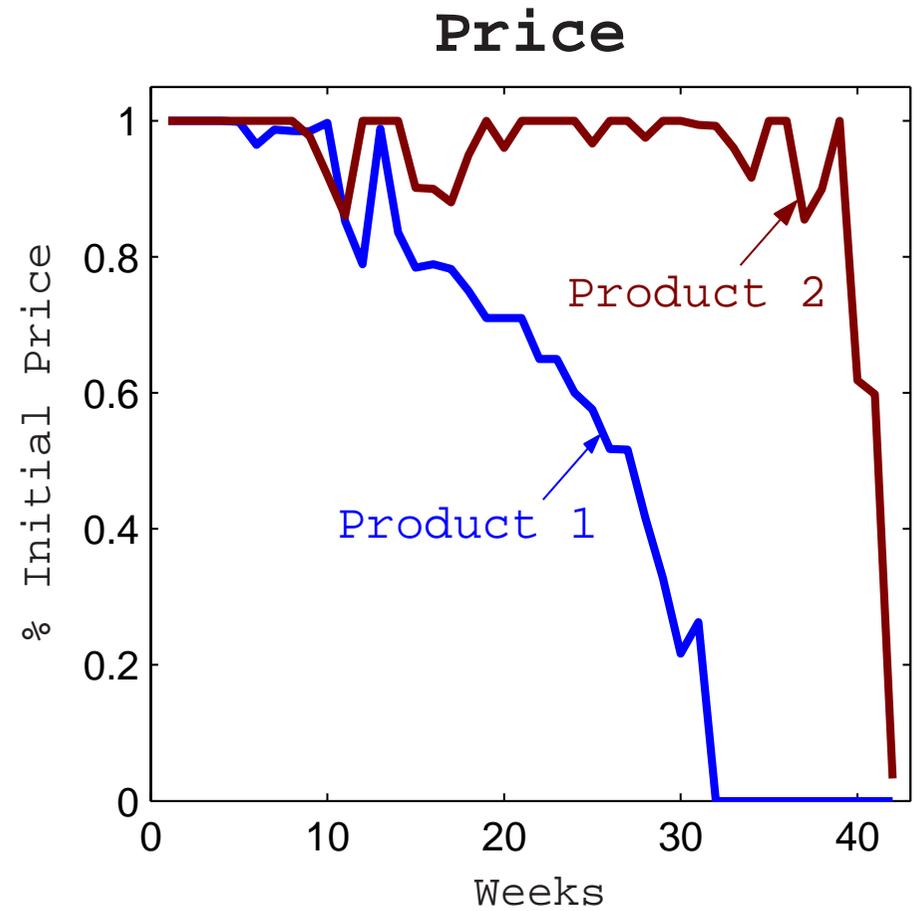
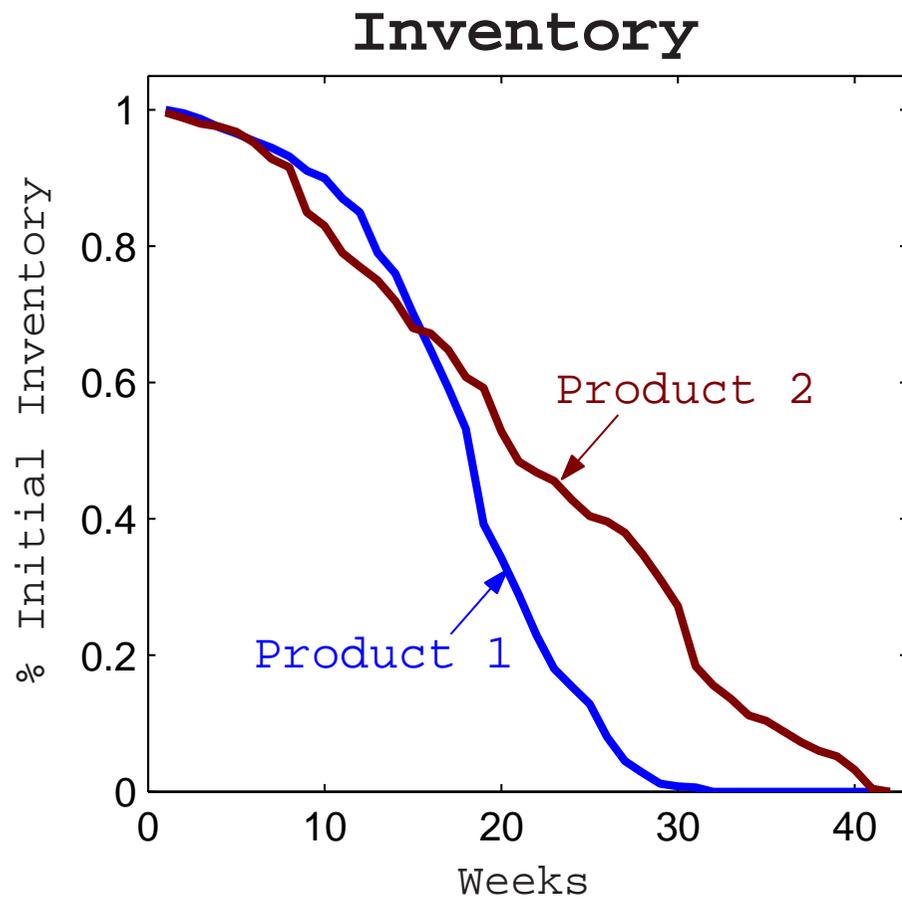
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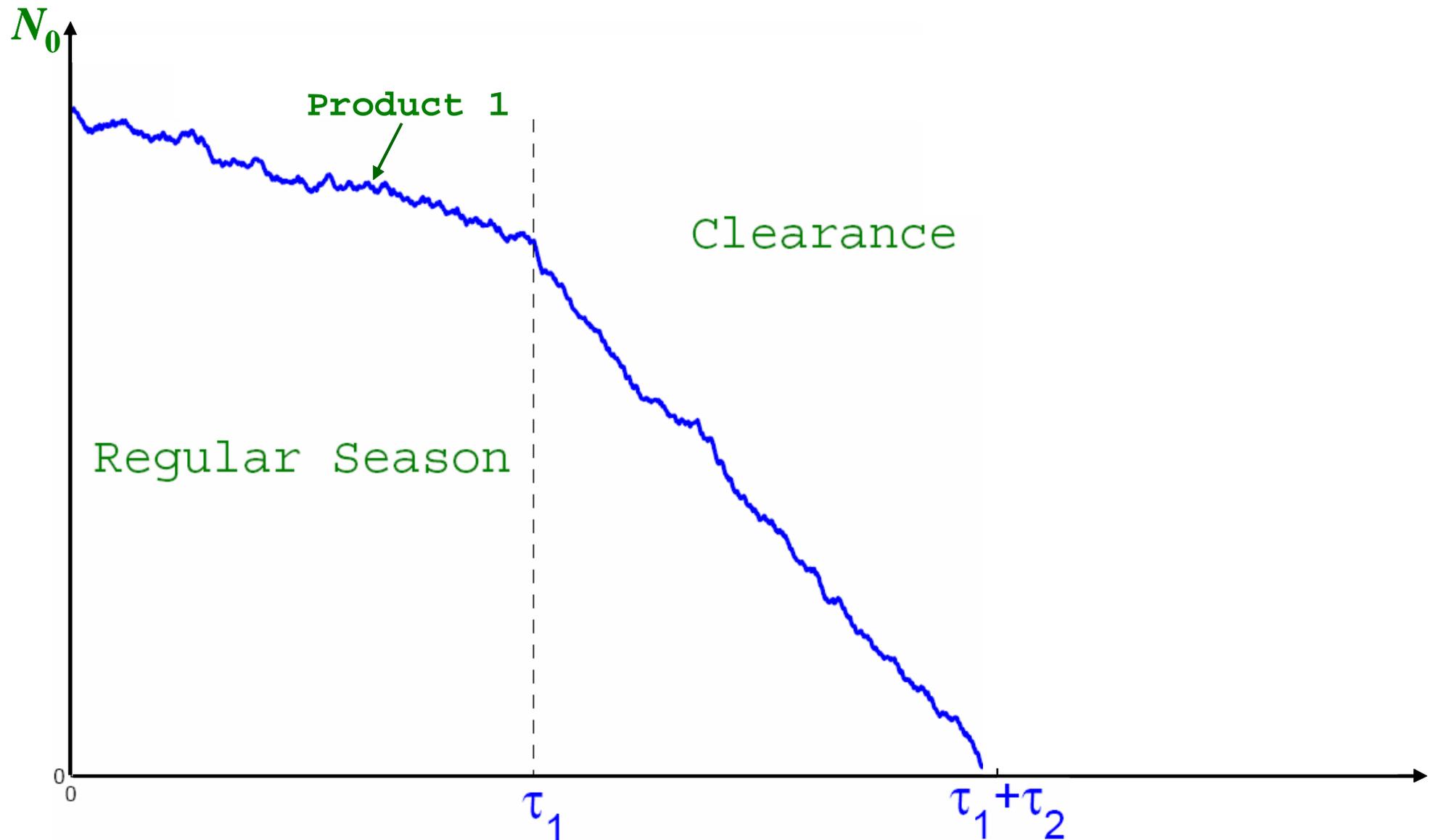
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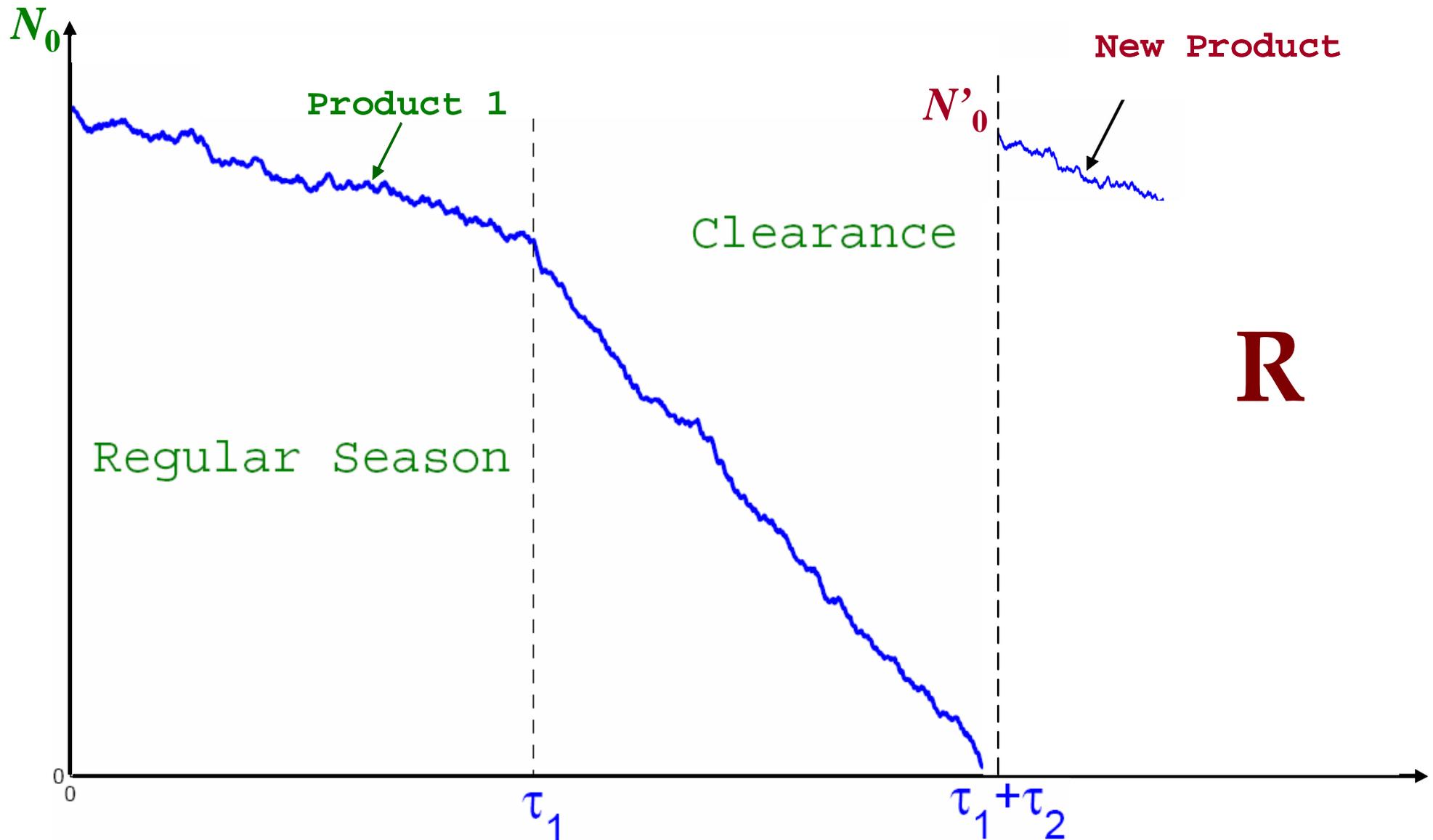
# Motivation



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# Motivation

- For many retail operations “capacity” is measured by store/shelf space.

- A key performance measure in the industry is

Average Sales per Square Foot per Unit Time.

- Trade-off between short-term benefits and the opportunity cost of assets.

Margin vs. Rotation.

- As opposed to the airline or hospitality industries, selling horizons are flexible.
- In general, most profitable/unprofitable products are new items for which there is little demand information.

# Outline

- ✓ Model Formulation.
- ✓ Perfect Demand Information.
- ✓ Incomplete Demand Information.
  - Inventory Clearance
  - Optimal Stopping (“outlet option”)
- ✓ Conclusion.

# Model Formulation

## I) STOCHASTIC SETTING:

- A probability space  $(\Omega, \mathcal{F}, \mathbb{P})$ .
- A standard Poisson process  $D(t)$  under  $\mathbb{P}$  and its filtration  $\mathcal{F}_t = \sigma(D(s) : 0 \leq s \leq t)$ .
- A collection  $\{\mathbb{P}_\alpha : \alpha > 0\}$  such that  $D(t)$  is a Poisson process with intensity  $\alpha$  under  $\mathbb{P}_\alpha$ .
- For a process  $f_t$ , we define  $I_f(t) := \int_0^t f_s ds$ .

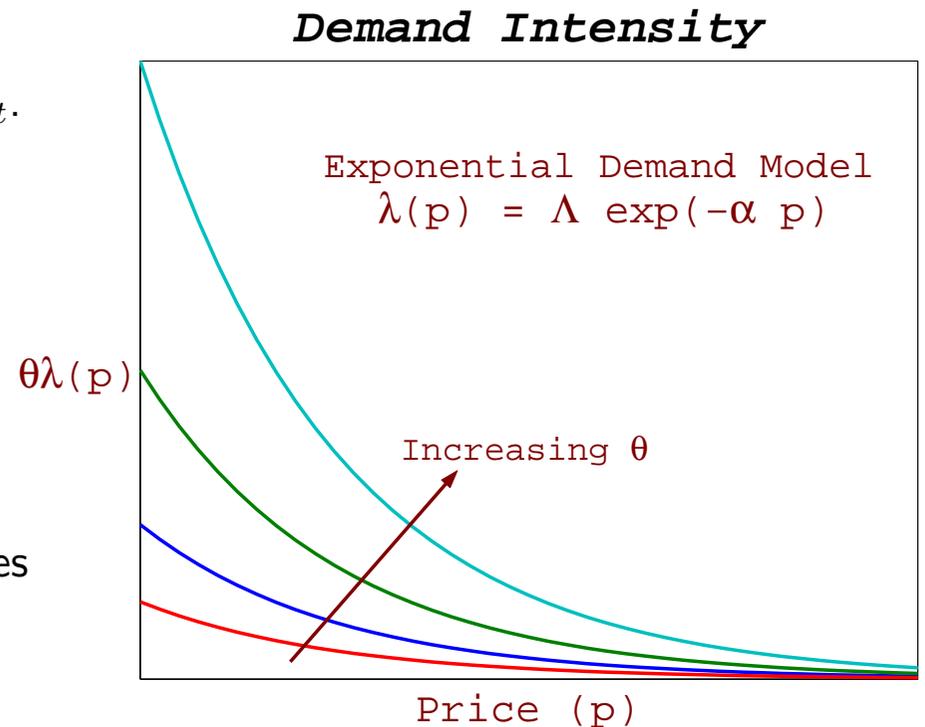
## II) DEMAND PROCESS:

- Pricing strategy, a nonnegative (adapted) process  $p_t$ .
- A price-sensitive **unscaled** demand intensity

$$\lambda_t := \lambda(p_t) \iff p_t = p(\lambda_t).$$

- A (possibly unknown) demand **scale factor**  $\theta > 0$ .
- Cumulative demand process  $D(I_\lambda(t))$  under  $\mathbb{P}_\theta$ .
- Select  $\lambda \in \mathcal{A}$  the set of admissible (adapted) policies

$$\lambda_t : \mathbb{R}_+ \rightarrow [0, \Lambda].$$



# Model Formulation

## III) REVENUES:

- Unscaled revenue rate  $c(\lambda) := \lambda p(\lambda)$ ,  $\lambda^* := \operatorname{argmax}_{\lambda \in [0, \Lambda]} \{c(\lambda)\}$ ,  $c^* := c(\lambda^*)$ .
- Terminal value (opportunity cost):  $R$                       Discount factor:  $r$
- Normalization:  $c^* = r R$ .

## IV) SELLING HORIZON:

- Inventory position:  $N_t = N_0 - D(I_\lambda(t))$ .
- $\tau_0 = \inf\{t \geq 0 : N_t = 0\}$ ,  $\mathcal{T} := \{\mathcal{F}_t - \text{stopping times } \tau \text{ such that } \tau \leq \tau_0\}$

## V) RETAILER'S PROBLEM:

$$\begin{aligned} & \max_{\lambda \in \mathcal{A}, \tau \in \mathcal{T}} && \mathbb{E}_\theta \left[ \int_0^\tau e^{-rt} p(\lambda_t) dD(I_\lambda(t)) + e^{-r\tau} R \right] \\ & \text{subject to} && N_t = N_0 - D(I_\lambda(t)). \end{aligned}$$

# Full Information

Suppose  $\theta > 0$  is known at  $t = 0$  and an **inventory clearance** strategy is used, i.e.,  $\tau = \tau_0$ .

Define the value function

$$W(n; \theta) = \max_{\lambda \in \mathcal{A}} \mathbb{E}_\theta \left[ \int_0^{\tau_0} e^{-rt} p(\lambda_t) dD(I_\lambda(t)) + e^{-r\tau} R \right]$$

subject to  $N_t = n - D(I_\lambda(t))$  and  $\tau_0 = \inf\{t \geq 0 : N_t = 0\}$ .

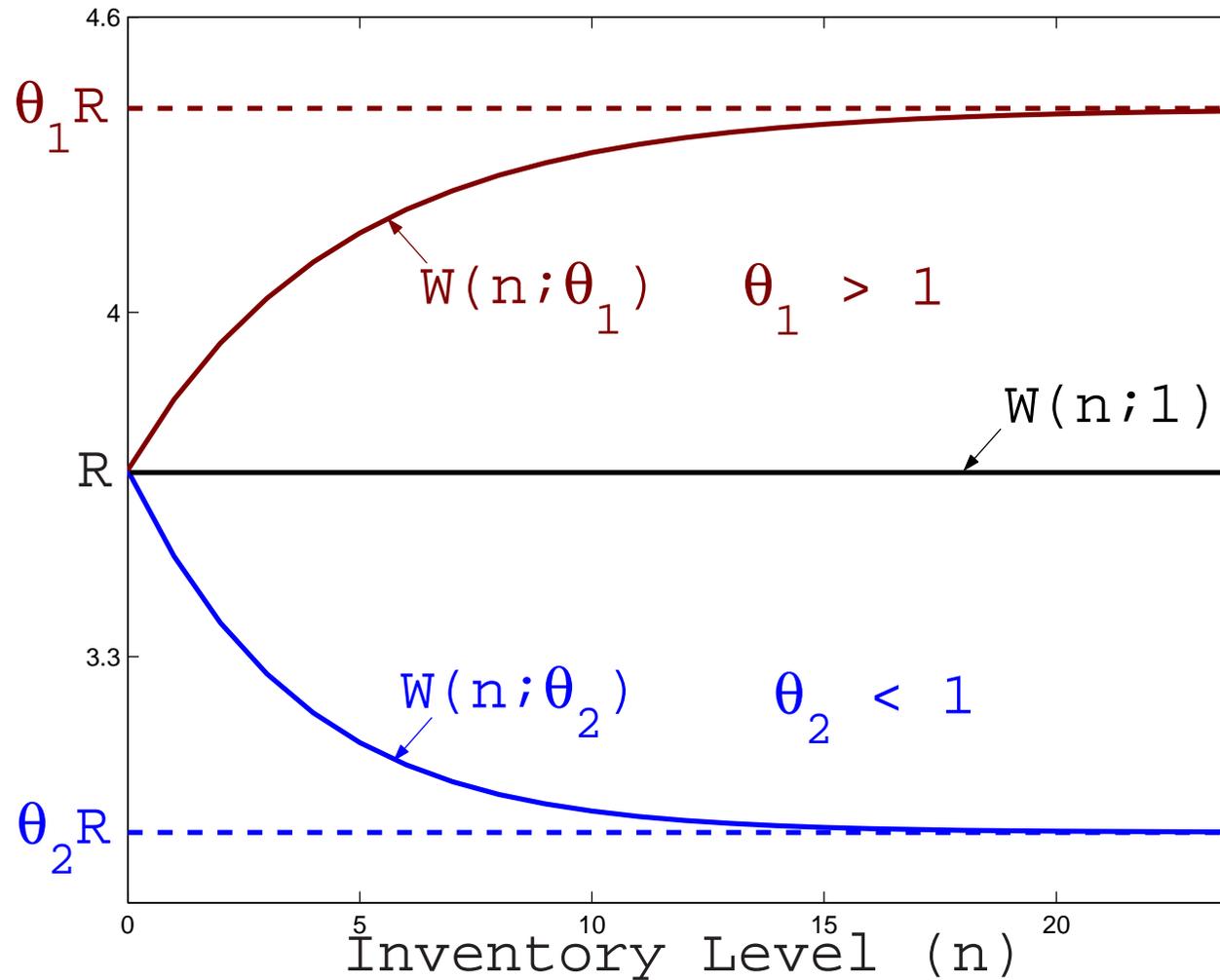
The solution satisfies the recursion  $\frac{r W(n; \theta)}{\theta} = \Psi(W(n-1; \theta) - W(n; \theta))$  and  $W(0; \theta) = R$ ,

$$\text{where } \Psi(z) \triangleq \max_{0 \leq \lambda \leq \Lambda} \{\lambda z + c(\lambda)\}.$$

**Proposition.** For every  $\theta > 0$  and  $R \geq 0$  there is a unique solution  $\{W(n) : n \in \mathbb{N}\}$ .

- If  $\theta \geq 1$  then the value function  $W$  is increasing and concave as a function of  $n$ .
- If  $\theta \leq 1$  then the value function  $W$  is decreasing and convex as a function of  $n$ .
- $\lim_{n \rightarrow \infty} W(n) = \theta R$ .

# Full Information



Value function for two values of  $\theta$  and an exponential demand rate  $\lambda(p) = \Lambda \exp(-\alpha p)$ .  
The data used is  $\Lambda = 10$ ,  $\alpha = 1$ ,  $r = 1$ ,  $\theta_1 = 1.2$ ,  $\theta_2 = 0.8$ ,  $R = \Lambda \exp(-1)/(\alpha r) \approx 3.68$ .

# Full Information

**Corollary.** Suppose  $c(\lambda)$  is strictly concave.

The optimal sales intensity satisfies:

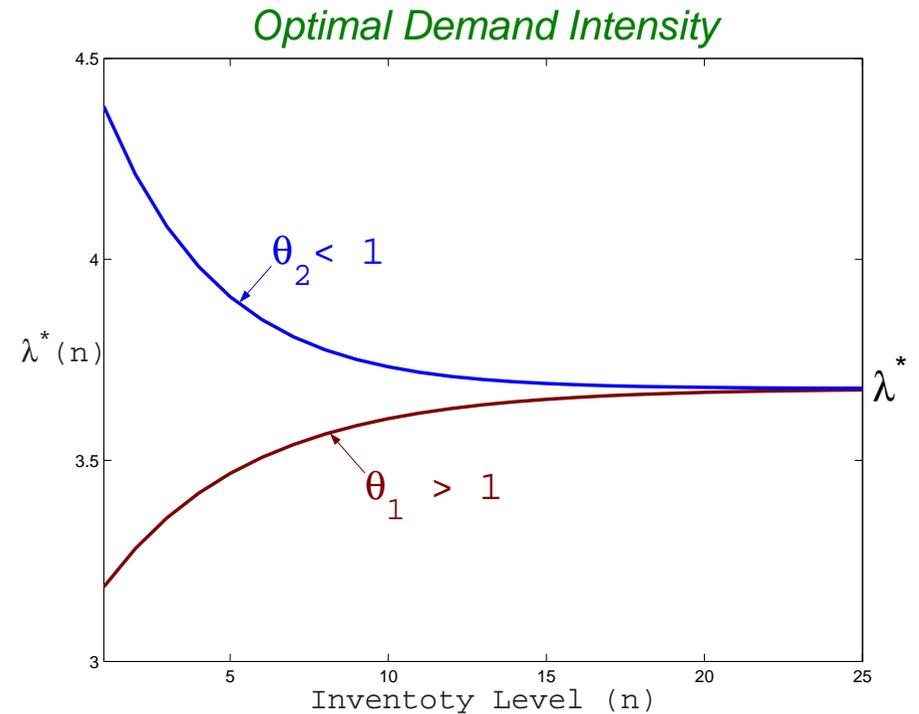
$$\lambda^*(n; \theta) = \operatorname{argmax}_{0 \leq \lambda \leq \Lambda} \{ \lambda (W(n-1; \theta) - W(n; \theta)) + c(\lambda) \}.$$

- If  $\theta \geq 1$  then  $\lambda^*(n; \theta) \uparrow n$ .

- If  $\theta \leq 1$  then  $\lambda^*(n; \theta) \downarrow n$ .

-  $\lambda^*(n; \theta) \downarrow \theta$ .

-  $\lim_{n \rightarrow \infty} \lambda^*(n, \theta) = \lambda^*$ .



Exponential Demand  $\lambda(p) = \Lambda \exp(-\alpha p)$ .  
 $\Lambda = 10, \alpha = r = 1, \theta_1 = 1.2, \theta_2 = 0.8, R = 3.68$ .

What about inventory turns (rotation)?

**Proposition.** Let  $s(n, \theta) \triangleq \theta \lambda^*(n, \theta)$  be the optimal sales rate for a given  $\theta$  and  $n$ .

$$\text{If } \frac{d}{d\lambda}(\lambda p'(\lambda)) \leq 0, \quad \text{then } s(n, \theta) \uparrow \theta.$$

# Full Information

## SUMMARY:

- A tractable dynamic pricing formulation for the inventory clearance model.
- $W(n; \theta)$  satisfies a simple recursion based on the Fenchel-Legendre transform of  $c(\lambda)$ .
- With full information products are divided in two groups:
  - High Demand Products with  $\theta \geq 1$ :  $W(n, \theta)$  and  $\lambda^*(n)$  increase with  $n$ .
  - Low Demand Products with  $\theta \leq 1$ :  $W(n, \theta)$  and  $\lambda^*(n)$  decrease with  $n$ .
- High Demand products are sold at a higher price and have a higher selling rate.
- If the retailer can stop selling the product at any time at no cost then:
  - If  $\theta < 1$  stop immediately ( $\tau = 0$ ).
  - If  $\theta > 1$  never stop ( $\tau = \tau_0$ ).
- In practice, a retailer rarely knows the value of  $\theta$  at  $t = 0$ !

# Incomplete Information: Inventory Clearance

## SETTING:

- The retailer does not know  $\theta$  at  $t = 0$  but knows  $\theta \in \{\theta_L, \theta_H\}$  with  $\theta_L \leq 1 \leq \theta_H$ .
- The retailer has a prior belief  $q \in (0, 1)$  that  $\theta = \theta_L$ .
- We introduce the probability measure  $\mathbb{P}_q = q \mathbb{P}_{\theta_L} + (1 - q) \mathbb{P}_{\theta_H}$ .
- We assume an inventory clearance model, *i.e.*,  $\tau = \tau_0$ .

## RETAILER'S BELIEFS:

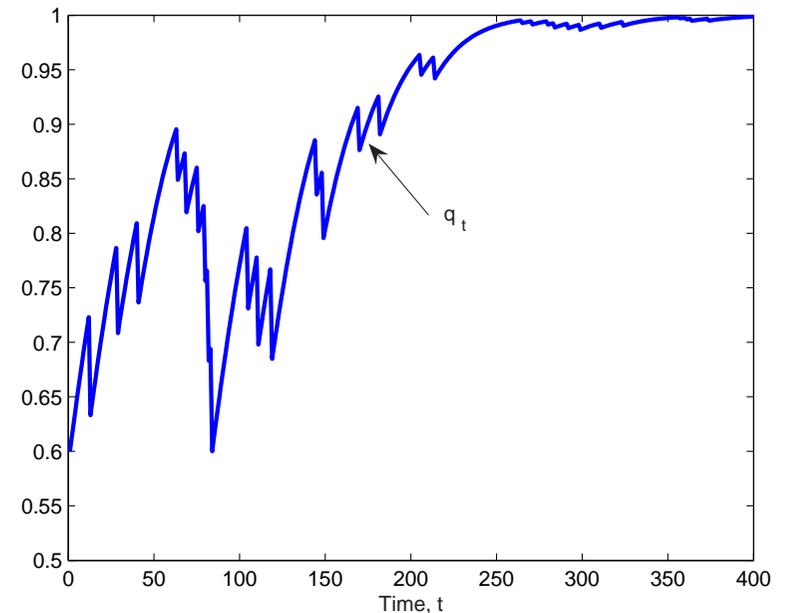
Define the belief process  $q_t := \mathbb{P}_q[\theta \mid \mathcal{F}_t]$ .

**Proposition.**  $q_t$  is a  $\mathbb{P}_q$ -martingale that satisfies the SDE

$$dq_t = -\eta(q_{t-}) [dD_t - \lambda_t \bar{\theta}(q_{t-}) dt],$$

$$\text{where } \bar{\theta}(q) := \theta_L q + \theta_H (1 - q)$$

$$\text{and } \eta(q) := \frac{q(1-q)(\theta_H - \theta_L)}{\theta_L q + \theta_H (1 - q)}.$$



# Incomplete Information: Inventory Clearance

## RETAILER'S OPTIMIZATION:

$$V(N_0, q) = \sup_{\lambda \in \mathcal{A}} \mathbb{E}_q \left[ \int_0^{\tau_0} e^{-rt} p(\lambda_t) dD(I_\lambda(s)) + e^{-r\tau_0} R \right]$$

subject to

$$N_t = N_0 - \int_0^t dD(I_\lambda(s)),$$

$$dq_t = -\eta(q_{t-}) [dD_t - \lambda_t \bar{\theta}(q_{t-}) dt], \quad q_0 = q,$$

$$\tau_0 = \inf\{t \geq 0 : N_t = 0\}.$$

## HJB EQUATION:

$$rV(n, q) = \max_{0 \leq \lambda \leq \Lambda} \left[ \lambda \bar{\theta}(q) [V(n-1, q - \eta(q)) - V(n, q) + \eta(q) V_q(n, q)] + \bar{\theta}(q) c(\lambda) \right],$$

with boundary condition  $V(0, q) = R$ ,  $V(n, 0) = W(n; \theta_H)$ , and  $V(n, 1) = W(n; \theta_L)$ .

## RECURSIVE SOLUTION:

$$V(0, q) = R, \quad V(n, q) + \Phi \left( \frac{r V(n, q)}{\bar{\theta}(q)} \right) - \eta(q) V_q(n, q) = V(n-1, q - \eta(q)).$$

# Incomplete Information: Inventory Clearance

## Proposition.

- ) The value function  $V(n, q)$  is
  - a) monotonically decreasing and convex in  $q$ ,
  - b) bounded by

$$W(n; \theta_L) \leq V(n, q) \leq W(n; \theta_H), \quad \text{and}$$

- c) uniformly convergent as  $n \uparrow \infty$ ,

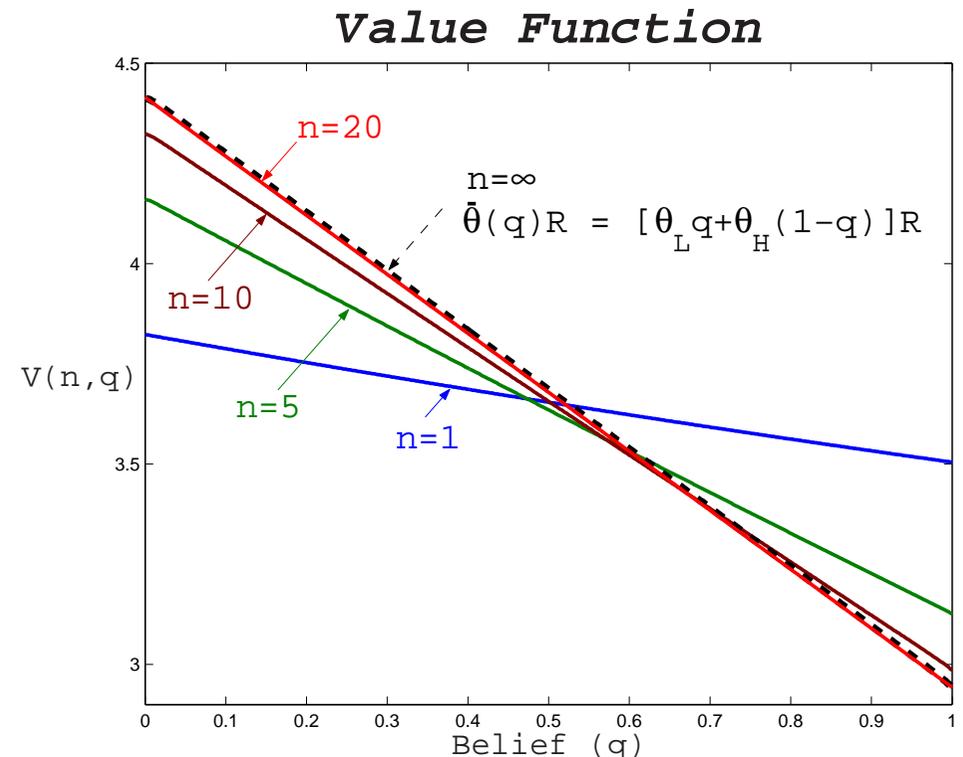
$$V(n, q) \xrightarrow{n \rightarrow \infty} R \bar{\theta}(q), \quad \text{uniformly in } q.$$

- ) The optimal demand intensity satisfies

$$\lim_{n \rightarrow \infty} \lambda^*(n, q) = \lambda^*.$$

## Conjecture:

The optimal sales rate  $\bar{\theta}(q) \lambda^*(n, q) \downarrow q$ .



# Incomplete Information: Inventory Clearance

ASYMPTOTIC APPROXIMATION: Since

$$\lim_{n \rightarrow \infty} V(n, q) = R \bar{\theta}(q) = \lim_{n \rightarrow \infty} \{q W(n, \theta_L) + (1 - q) W(n, \theta_H)\},$$

we propose the following approximation for  $V(n, q)$

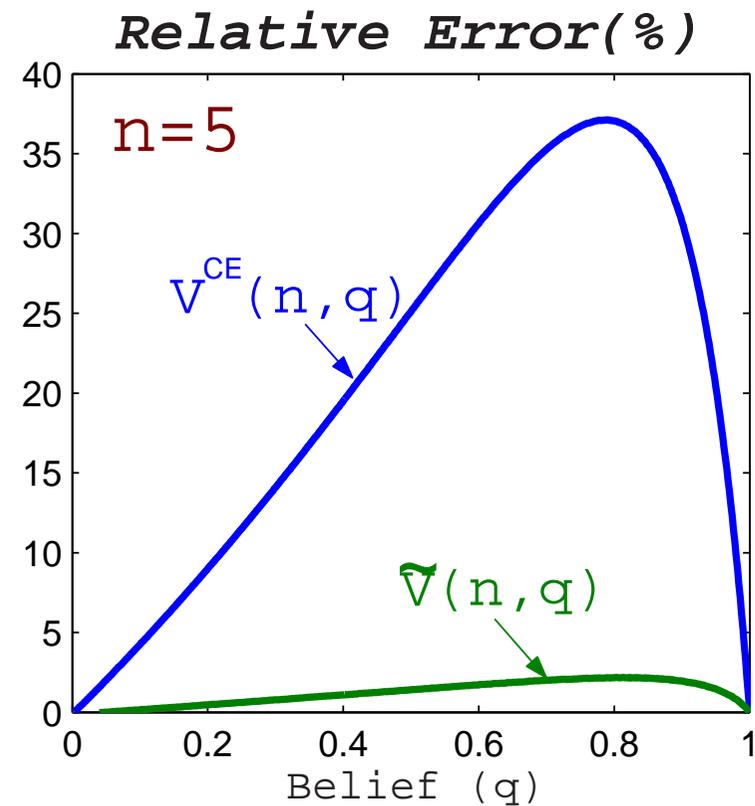
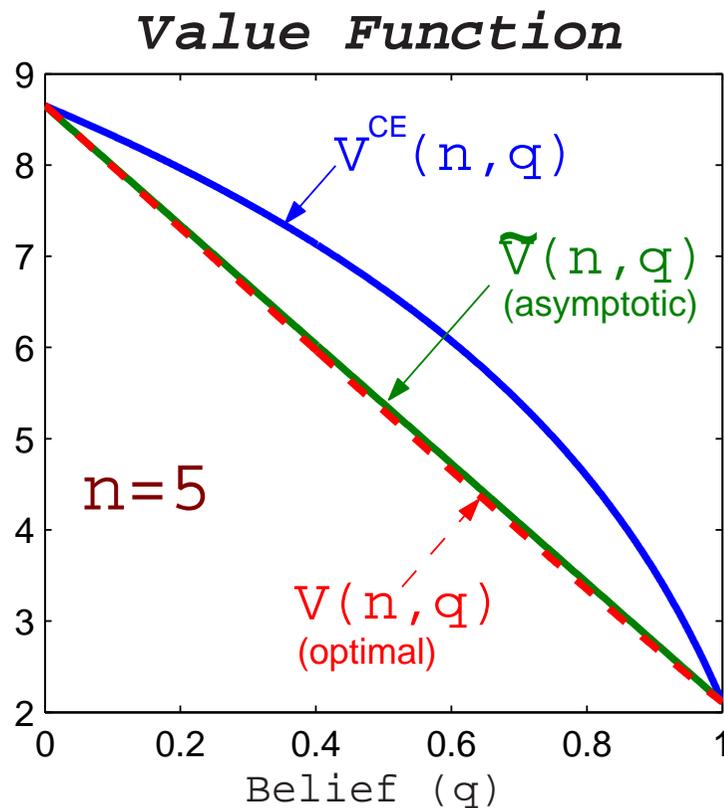
$$\tilde{V}(n, q) := q W(n, \theta_L) + (1 - q) W(n, \theta_H).$$

SOME PROPERTIES OF  $\tilde{V}(n, q)$ :

- Linear approximation easy to compute.
- Asymptotically optimal as  $n \rightarrow \infty$ .
- Asymptotically optimal as  $q \rightarrow 0^+$  or  $q \rightarrow 1^-$ .
- $\tilde{V}(n, q) = \mathbb{E}_q[W(n, \theta)] \neq W(n, \mathbb{E}_q[\theta]) =: V^{\text{CE}}(n, q) = \text{Certainty Equivalent.}$

# Incomplete Information: Inventory Clearance

$$\text{Relative Error (\%)} := \frac{V^{\text{approx}}(n, q) - V(n, q)}{V(n, q)} \times 100\%.$$



Exponential Demand  $\lambda(p) = \Lambda \exp(-\alpha p)$ :

Inventory = 5,  $\Lambda = 10$ ,  $\alpha = r = 1$ ,  $\theta_H = 5.0$ ,  $\theta_L = 0.5$ .

# Incomplete Information: Inventory Clearance

For any approximation  $V^{\text{approx}}(n, q)$ , define the corresponding demand intensity using the HJB

$$\lambda^{\text{approx}}(n, q) := \arg \max_{0 \leq \lambda \leq \Lambda} [\lambda \bar{\theta}(q) [V^{\text{approx}}(n-1, q-\eta(q)) - V^{\text{approx}}(n, q)] + \lambda \kappa(q) V_q^{\text{approx}}(n, q) + \bar{\theta}(q) c(\lambda)].$$

$$\text{Relative Price Error (\%)} := \frac{p(\lambda^{\text{approx}}) - p(\lambda^*)}{p(\lambda^*)} \times 100\%.$$

ASYMPTOTIC APPROXIMATION (%)

	Inventory ( $n$ )				
q	1	5	10	25	100
0.0	0.0	0.0	0.0	0.0	0.0
0.2	2.7	-0.2	-0.3	-0.6	-0.5
0.4	6.9	<b>0.8</b>	<b>-0.6</b>	<b>-0.9</b>	-0.7
0.6	12.5	<b>2.4</b>	<b>-0.2</b>	<b>-0.7</b>	-1.0
0.8	19.4	<b>3.3</b>	<b>0.1</b>	<b>-0.4</b>	-0.6
1.0	0.0	0.0	0.0	0.0	0.0

CERTAINTY EQUIVALENT (%)

	Inventory ( $n$ )				
q	1	5	10	25	100
0.0	0.0	0.0	0.0	0.0	0.0
0.2	5.3	2.6	2.7	2.4	-0.4
0.4	14.4	<b>11.6</b>	<b>12.0</b>	<b>10.1</b>	-0.5
0.6	29.9	<b>28.2</b>	<b>28.0</b>	<b>17.6</b>	-1.0
0.8	54.6	<b>46.2</b>	<b>37.4</b>	<b>11.1</b>	-0.7
1.0	0.0	0.0	0.0	0.0	0.0

Relative price error for the exponential demand model  $\lambda(p) = \Lambda \exp(-\alpha p)$ , with  $\Lambda = 20$  and  $\alpha = 1$ .

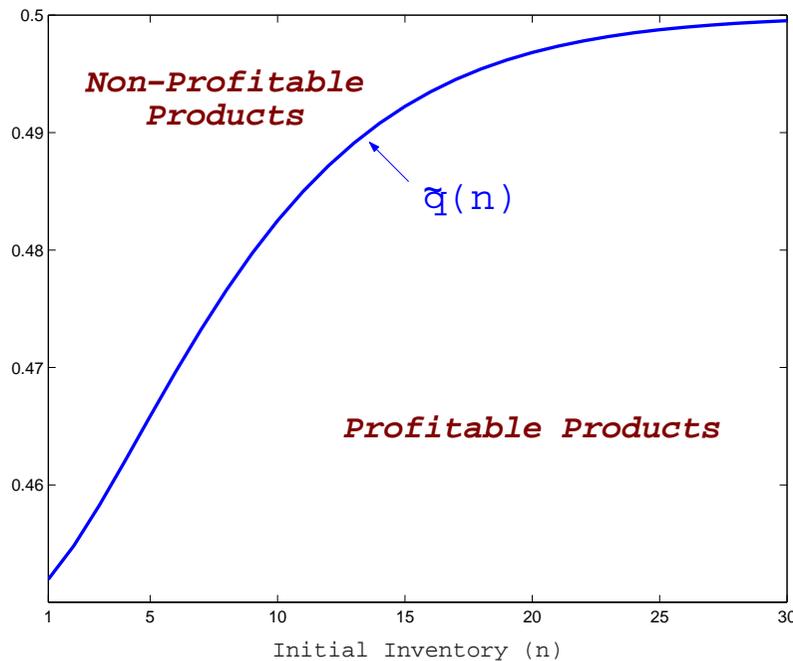
# Incomplete Information: Inventory Clearance

When should the retailer engage in selling a given product?

When  $V(n, q) \geq R$ .

Using the asymptotic approximation  $\tilde{V}(n, q)$ , this is equivalent to

$$q \leq \tilde{q}(n) := \frac{W(n; \theta_H) - R}{[W(n; \theta_H) - R] + [R - W(n; \theta_L)]}.$$



Exponential demand rate  $\lambda(p) = \Lambda \exp(-\alpha p)$ .  
 Data:  $\Lambda = 10$ ,  $\alpha = 1$ ,  $r = 1$ ,  $\theta_H = 1.2$ ,  $\theta_L = 0.8$ .

$$\tilde{q}(n) \rightarrow \tilde{q}_\infty := \frac{\theta_H - 1}{\theta_H - \theta_L}, \text{ as } n \rightarrow \infty.$$

# Incomplete Information: Inventory Clearance

## SUMMARY:

- Uncertainty in market size ( $\theta$ ) is captured by a new state variable  $q_t$  (a jump process).
- $V(n, q)$  can be computed using a recursive sequence of ODEs.
- Asymptotic approximation  $\tilde{V}(n, q) := \mathbb{E}_q[W(n, \theta)]$  performs quite well.
  - Linear approximation easy to compute.
  - Value function:  $V(n, q) \approx \tilde{V}(n, q)$ .
  - Pricing strategy:  $p^*(n, q) \approx \tilde{p}(n, q)$ .
- Products are divided in two groups as a function of  $(n, q)$ :
  - Profitable Products with  $q < \tilde{q}(n)$  and
  - Non-profitable Products with  $q > \tilde{q}(n)$ .
- The threshold  $\tilde{q}(n)$  increases with  $n$ , that is, the retailer is willing to take more risk for larger orders.

# Incomplete Information: Optimal Stopping

## SETTING:

- Retailer does not know  $\theta$  at  $t = 0$  but knows  $\theta \in \{\theta_L, \theta_H\}$  with  $\theta_L \leq 1 \leq \theta_H$ .
- Retailer has the option of removing the product at any time, “Outlet Option”.

## RETAILER'S OPTIMIZATION:

$$U(N_0, q) = \max_{\lambda \in \mathcal{A}, \tau \in \mathcal{T}} \mathbb{E}_q \left[ \int_0^\tau e^{-rt} p(\lambda_t) dD(I_\lambda(t)) + e^{-r\tau} R \right]$$

subject to

$$N_t = N_0 - D(I_\lambda(t)),$$

$$dq_t = -\eta(q_{t-}) [dD(I_\lambda(t)) - \lambda_t \bar{\theta}(q_{t-}) dt], \quad q_0 = q.$$

## OPTIMALITY CONDITIONS:

$$\begin{cases} U(n, q) + \Phi\left(\frac{rU(n, q)}{\bar{\theta}(q)}\right) - \eta(q) U_q(n, q) = U(n-1, q - \eta(q)) & \text{if } U \geq R \\ U(n, q) + \Phi\left(\frac{rU(n, q)}{\bar{\theta}(q)}\right) - \eta(q) U_q(n, q) \leq U(n-1, q - \eta(q)) & \text{if } U = R. \end{cases}$$

# Incomplete Information: Optimal Stopping

## Proposition.

- a) There is a unique continuously differentiable solution  $U(n, \cdot)$  defined on  $[0, 1]$  so that  $U(n, q) > R$  on  $[0, q_n^*)$  and  $U(n, q) = R$  on  $[q_n^*, 1]$ , where  $q_n^*$  is the unique solution of

$$R + \Phi\left(\frac{rR}{\bar{\theta}(q)}\right) = U(n-1, q - \eta(q)).$$

- b)  $q_n^*$  is increasing in  $n$  and satisfies

$$\frac{\theta_H - 1}{\theta_H - \theta_L} \leq q_n^* \xrightarrow{n \rightarrow \infty} q_\infty^* \leq \text{Root} \left\{ \Phi\left(\frac{rR}{\bar{\theta}(q)}\right) = \frac{\eta(q)}{q} (\theta_H - 1) R \right\} < 1.$$

- c) The value function  $U(n, q)$

- Is decreasing and convex in  $q$  on  $[0, 1]$
- Increases in  $n$  for all  $q \in [0, 1]$  and satisfies

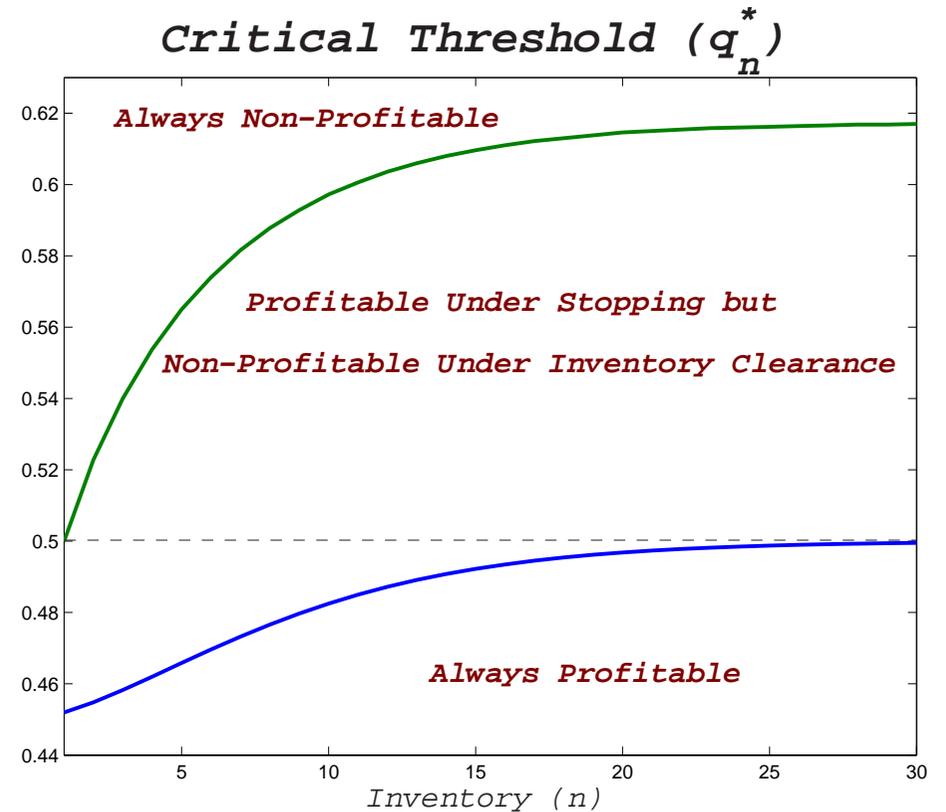
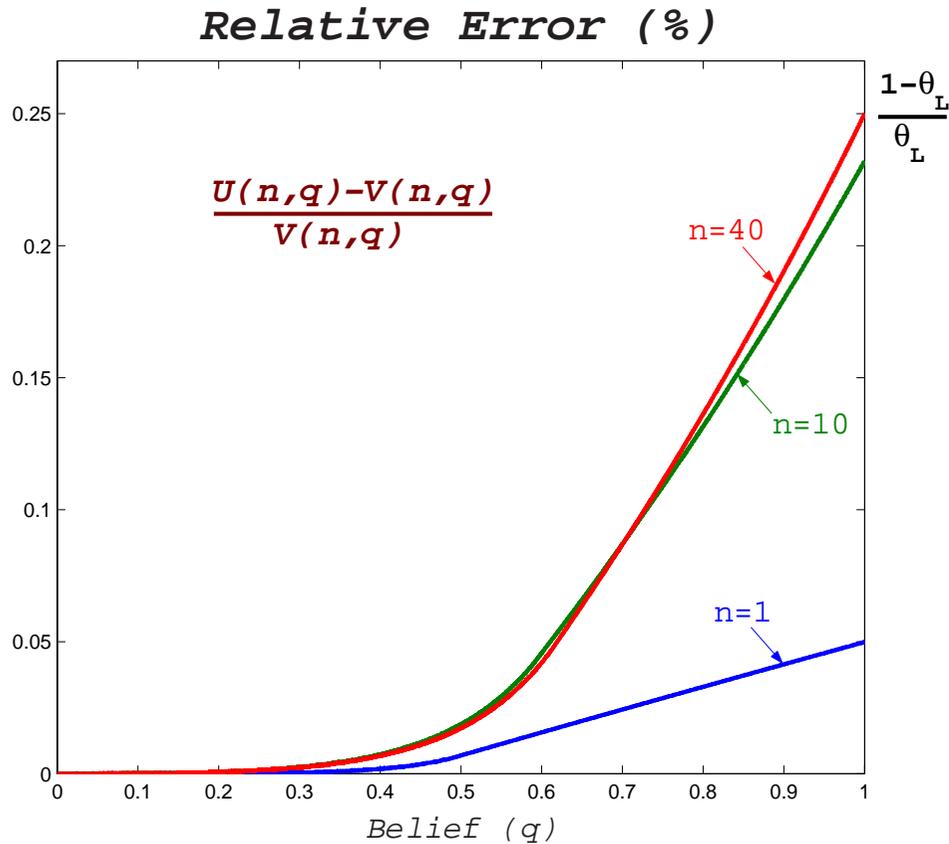
$$\max\{R, V(n, q)\} \leq U(n, q) \leq \max\{R, m(q)\} \quad \text{for all } q \in [0, 1],$$

where

$$m(q) := W(n, \theta_H) - \frac{(W(n, \theta_H) - R)}{q_n^*} q.$$

- Converges uniformly (in  $q$ ) to a continuously differentiable function,  $U_\infty(q)$ .

# Incomplete Information: Optimal Stopping



Exponential demand rate  $\lambda(p) = \Lambda \exp(-\alpha p)$ .

Data:  $\Lambda = 10$ ,  $\alpha = 1$ ,  $r = 1$ ,  $\theta_H = 1.2$ ,  $\theta_L = 0.8$ .

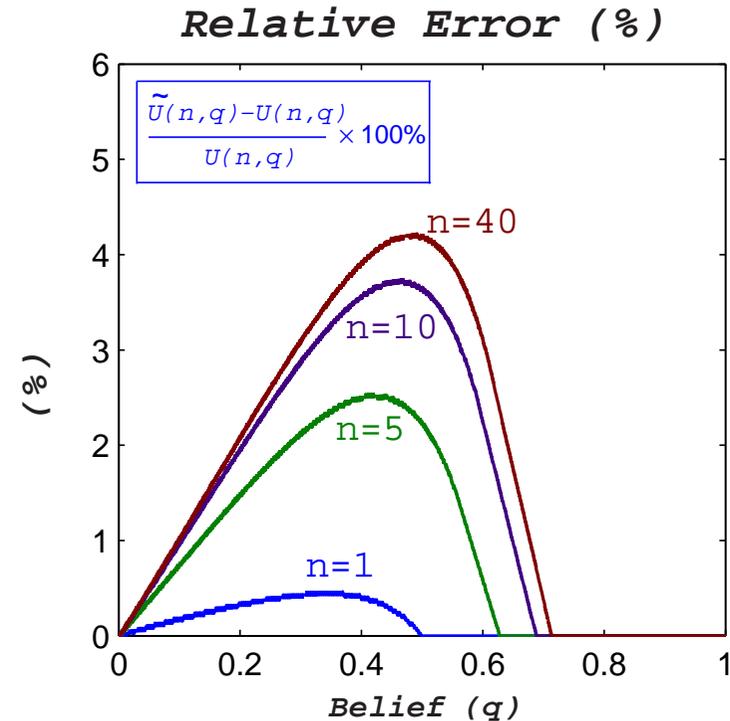
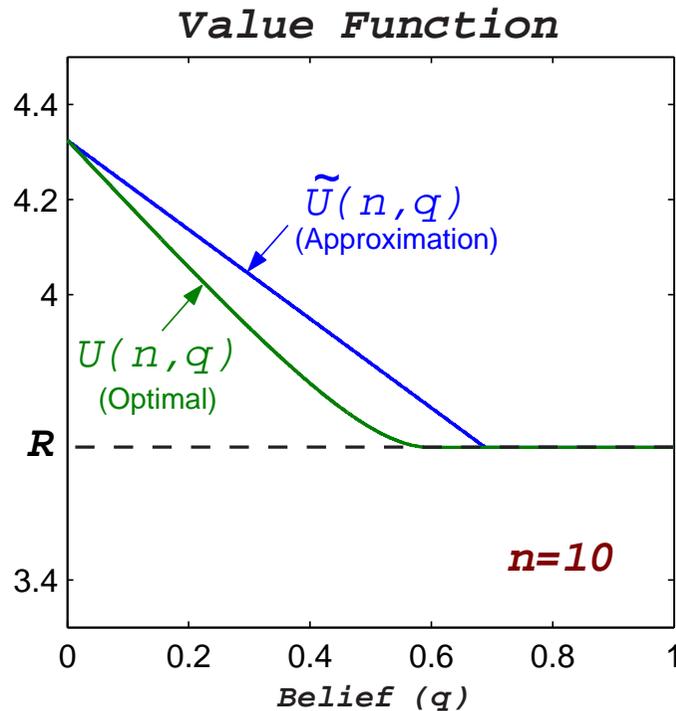
# Incomplete Information: Optimal Stopping

## APPROXIMATION:

$$\tilde{U}(n, q) := \max\left\{R, W(n, \theta_H) - \frac{(W(n, \theta_H) - R)}{\tilde{q}_n} q\right\}$$

where  $\tilde{q}_n$  is the unique solution of

$$R + \Phi\left(\frac{rR}{\bar{\theta}(q)}\right) = \tilde{U}(n-1, q - \eta(q)).$$



Exponential demand rate  $\lambda(p) = \Lambda \exp(-\alpha p)$ .

Data:  $\Lambda = 10$ ,  $\alpha = 1$ ,  $r = 1$ ,  $\theta_H = 1.2$ ,  $\theta_L = 0.8$ .

# Incomplete Information: Optimal Stopping

## SUMMARY:

- $U(n, q)$  can be computed using a recursive sequence of ODEs with free-boundary conditions.
- For every  $n$  there is a critical belief  $q_n^*$  above which it is optimal to stop.
- Again, the sequence  $q_n^*$  is increasing with  $n$ , that is, the retailer is willing to take more risk for larger orders.
- The sequence  $q_n^*$  is bounded by

$$\frac{\theta_H - 1}{\theta_H - \theta_L} \leq q_n^* \leq \hat{q} := \text{Root} \left\{ \Phi \left( \frac{r R}{\bar{\theta}(q)} \right) = \frac{\eta(q)}{q} (\theta_H - 1) R \right\}$$

- The “outlet option” increases significantly the expected profits and the range of products  $(n, q)$  that are *profitable*.

$$0 \leq U(n, q) - V(n, q) \leq (1 - \theta_L)^+ R.$$

- A simple piece-wise linear approximation works well.

$$\tilde{U}(n, q) := \max \left\{ R, W(n, \theta_H) - \frac{(W(n, \theta_H) - R)}{\tilde{q}_n} q \right\}$$

# Concluding Remarks

- A simple dynamic pricing model for a retailer selling non-perishable products.
- Captures two common sources of uncertainty:
  - Market size measured by  $\theta \in \{\theta_H, \theta_L\}$ .
  - Stochastic arrival process of price sensitive customers.
- Analysis gets simpler using the Fenchel-Legendre transform of  $c(\lambda)$  and its properties.
- We propose a simple approximation (linear and piecewise linear) for the value function and corresponding pricing policy.
- Some properties of the optimal solution are:
  - Value functions  $V(n, q)$  and  $U(n, q)$  are decreasing and convex in  $q$ .
  - The retailer is willing to take more risk ( $\uparrow q$ ) for higher orders ( $\uparrow n$ ).
  - The optimal demand intensity  $\lambda^*(n, q) \uparrow q$  and the optimal sales rate  $\bar{\theta}(q) \lambda^*(n, q) \downarrow q$ .
- **Extension:**  $R(n) = R + \nu n - K \mathbb{1}(n > 0)$ .