



Reconnect '04


Experimental Algorithmics with a Focus on Branch and Bound for Discrete Optimization Problems

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Discrete (Combinatorial) Optimization


Informally, a Combinatorial Optimization problem is the minimization (maximization) of an **objective function** over a **finite** set of **feasible** solutions.

Examples from everyday life:

- Scheduling family activities (just finding a feasible solution)
 - Vehicle routing
- Packing a moving van or dishwasher
- Budgeting

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Complexity


Most of the problems we really want to solve are NP-complete: Any algorithm that computes the **optimal** solution for **all** input instances is likely to take an unacceptable amount of time in the **worst case**.

But we still need to solve them in practice:

- Approximation
- Restricted input instances
- Average Case

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Solving Combinatorial Optimization Problems in Practice

Goal: Do the best we can in solving a **particular instance**.

- That's what people really want to solve
- Structural insight from optimal solutions to small problems leads to
 - Better approximation algorithms (theory)
 - Better heuristics (practice)
 - Better solution of bigger problems


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Experimental Algorithmics

- Rigorous computer experimentation
 - How well does an algorithm/implementation perform?
 - Time, Approximation quality
 - Algorithm comparisons
- Algorithmic questions arising during development of robust, efficient software
 - Cache issues
 - Generation of synthetic data
- Algorithm Engineering

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Branch and Bound


Branch and Bound is an intelligent (enumerative) search procedure for discrete optimization problems.


$$\min_{x \in X} f(x)$$

Three required (problem-specific) procedures:


- Compute a lower bound (assuming min) $b(X)$

$$\forall x \in X, b(X) \leq f(x)$$
- Split a feasible region (e.g. over parameter/decision space)
- Find a candidate solution $x \in X$
 - Can fail
 - Require that it recognizes feasibility if X has only one point


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Branching (Splitting)

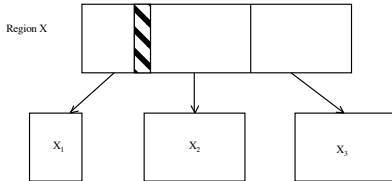


- Usually partitions the feasible region (or better)


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Branching (Splitting)



- Usually partitions the feasible region (or better)

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Branch and Bound

Root Problem = original

fathomed

infeasible

- Recursively divide feasible region, prune search when no optimal solution can be in the region.

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Terminology

- The best feasible solution found so far is called the **incumbent**.
- A node awaiting bounding or splitting (not done) is called **active**.
- The set of active nodes is sometimes called the **frontier**.

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Branch and Bound Solution Quality

- At any point in the computation, we have a global lower bound: the lowest lower bound among all open (active) nodes.
- We can stop at any point and compute an instance-specific approximation bound: value of incumbent/global bound
- If B&B runs to completion we have
 - Found an optimal solution
 - Proven that this solution is optimal

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Example: Binary Knapsack

Given set of objects $1..n$
 Each has weight/size w_i , and value v_i

This thief wants the most valuable set that fits into a knapsack of size W

Find a subset of objects S that
 Maximizes $V(S) = \sum_{i \in S} v_i$
 Subject to $W(S) = \sum_{i \in S} w_i \leq W$

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Binary Knapsack - Upper Bound (and Candidate)

- Bound with greedy strategy
 - Maximize value/unit weight
- Sort objects by "bang for the buck": $\frac{v_i}{w_i}$
 - Add objects in order
 - Split last object

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Greedy Isn't optimal

- Knapsack is weakly NP-complete

	w	v	v/w
A	6	12	2
B	5	9	9/5
C	4	4	1

- For $W = 9$, {B,C} is optimal

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Knapsack - Partitioning

Represent a decision region with three subsets: IN, OUT, UNDECIDED
 Upper bound and candidate: (value of) greedy on UNDECIDED with new bound $W - W(IN)$.
 Branch object (i above): The one partially-included item in the greedy solution

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Branch and Bound

- Recursively divide feasible region, prune search when no optimal solution can be in the region.

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Representation - Space issues

Two Lists (IN and OUT)
UNDECIDED is implied

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Representation Issue - Space

- 3-valued vector: IOU
- Save difference with parent
- Issue: length of chains vs. space cost of replication

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This course covers

- Practical ways to solve combinatorial optimization problems with branch and bound
- Integer Programming
 - Natural general way to express any NP-complete problem (most useful combinatorial optimization problems)
 - Theoretical foundations as relevant to practice
 - Software tools
- Experimental Algorithmics
 - Answering “so what?” for your experimental study

We'll see example applications, but they're the tip of the iceberg

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