

## Reconnect '04 Using PICO

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## AMPL Model for 1 | prec | $\sum w_j C_j$

- Declare/compute parameters, sets

```
param n > 0 integer;          # Number of job
set Jobs := 0..n-1;
param p {Jobs} >= 0 integer;  # Job Duration p(j)
param w {Jobs} >= 0 integer;  # Job weight (for objective) w(j)
# (predecessor, successor) pairs
set Precedence within Jobs cross Jobs;
# The makespan of the schedule (and latest finish time)
param T := sum {j in Jobs} p[j];
```

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## AMPL Model for 1 | prec | $\sum w_j C_j$

- Declare variables

```
var x {j in Jobs, t in p[j]..T} binary;
```

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## AMPL Model for 1 | prec | $\sum w_j C_j$

minimize WCT:

```
sum {j in Jobs, t in p[j]..T} w[j] * t * x[j,t];
```

- subject to TaskDone {j in Jobs}:  
sum {t in p[j]..T} x[j,t] = 1;
- subject to UnitWork {tu in 1..T}:  
sum {j in Jobs, t in tu..tu+p[j]-1 inter p[j]..T} x[j,t] <= 1;
- subject to PrecConstraint {j in Jobs, k in Jobs,  
tu in p[j]..T-p[k]: (j,k) in Precedence}:  
sum {t in p[j]..tu} x[j,t] >= sum {t in p[j]+p[k]..tu + p[k]} x[k,t];

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### AMPL Model for 1 | prec | $\sum w_j C_j$

minimize WCT:

$$\text{sum}\{j \text{ in Jobs}, t \text{ in } p[j]..T\} w[j] * t * x[j,t];$$

Corresponds to:

$$\min \sum_{j=1}^n \sum_{t=p_j}^T w_j t x_{jt}$$

subject to TaskDone {j in Jobs}:

$$\text{sum}\{t \text{ in } p[j]..T\} x[j,t] = 1;$$

Corresponds to:

$$\sum_{t=p_j}^T x_{jt} = 1 \quad \forall j$$

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### AMPL Model for 1 | prec | $\sum w_j C_j$

subject to UnitWork {tu in 1..T}:

$$\text{sum}\{j \text{ in Jobs}, t \text{ in } tu..tu+p[j]-1 \text{ inter } p[j]..T\} x[j,t] \leq 1;$$

Corresponds to:

$$\sum_{j=1}^n \sum_{u=t}^{t+p_j-1} x_{ju} \leq 1 \quad t = 1, \dots, T = \sum_{j=1}^n p_j$$

subject to PrecConstraint {j in Jobs, k in Jobs,

tu in p[j]..T-p[k]: (j,k) in Precedence}:

$$\text{sum}\{t \text{ in } p[j]..tu\} x[j,t] \geq \text{sum}\{t \text{ in } p[j]+p[k]..tu + p[k]\} x[k,t];$$

Corresponds to:

$$\sum_{u=1}^t x_{ju} - \sum_{u=1}^{t+p_k} x_{ku} \geq 0 \quad \forall J_i < J_k, t = 1, \dots, T - p_k$$

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### More Complex Sets to Reduce IP size

# Transitive closure of precedence relations.

set step{s in 1..ceil((n-1)/2)} dimen 2 :=

if s==1

then Precedence

else step[s-1] union setof {k in Jobs, (i,k) in step[s-1], (k,j) in step[s-1]} (i,j);

set TransPrec := step[ceil((n-1)/2)];

# Earliest finish time

param EFT{j in Jobs} := p[j] + sum{k in Jobs: (k,j) in TransPrec} p[k];

# Latest finish time

param LFT{j in Jobs} := T - sum{k in Jobs: (j,k) in TransPrec} p[k];

# Possible finish times

set FinishWindow{j in Jobs} := EFT[j]..LFT[j];

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### Smaller, simpler formulation

var x {j in Jobs, t in FinishWindow[j]} binary;

minimize WCT:

$$\text{sum}\{j \text{ in Jobs}, t \text{ in FinishWindow[j]}\} w[j] * t * x[j,t];$$

subject to TaskDone {j in Jobs}:

$$\text{sum}\{t \text{ in FinishWindow[j]}\} x[j,t] = 1;$$

subject to UnitWork {tu in 1..T}:

$$\text{sum}\{j \text{ in Jobs}, t \text{ in } tu..tu+p[j]-1 \text{ inter FinishWindow[j]}\} x[j,t] \leq 1;$$

[Similar changes for precedence constraints]

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## AMPL Data Files - Straightforward

```

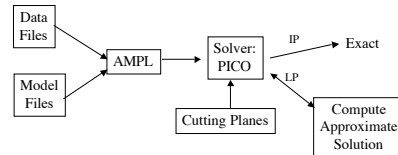
param n := 6;
param p :=
0 8
1 2
2 1
3 3
4 12
5 1;
set Precedence :=
0 2
1 2
2 3
4 5;

```

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## Use Sandia National Labs IP Solver: PICO



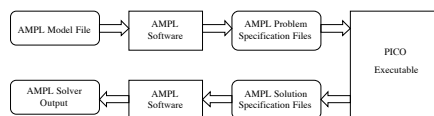
- Write cutting-plane and approximate-solution code using AMPL variables
- Mapping transparent

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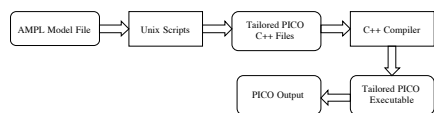


## AMPL-PICO Interface

### Standard AMPL interfaces



### Customized PICO Interface



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## Caveats: Status as of June 22, 2004

- The derived classes are for branch and bound only (no cuts yet; not stable; but they'll be there "soon")
- Only the serial version works as of this morning
  - This could change by tomorrow

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## Customized PICO Application

Goal: transparent access to AMPL variables from within problem-specific derivatives of core PICO application classes:

- MILP The “master” branching class
- MILPNode The “worker” subproblem class
- MILPProblem The problem definition class

- Creates a new namespace, so (as far as you’re concerned), these names are the same.
- Maps multidimensional ampl variables to PICO’s linear variable space
  - Foo(1,2)
  - Foo(bar(stuff),callFunc(moreStuff))

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## How PICO Learns AMPL Names

- Variable names, constraint names
  - Automatic (when the customized code is set up)
- Other parameters: list explicitly
  - You may never use some helper parameters from the ampl file

# PICO SYMBOL: p n w EFT LFT T

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## Example: Incumbent Heuristic for $1/\text{prec} \sum w_j C_j$

Computing midpoints by stepping through valid x in lexicographic order

```
DoubleVector frac(n());
IntVector time(n());
int curr_job = -1;
x_const_iterator curr = x_valid().begin();
x_const_iterator last = x_valid().end();
while (curr != last) {
    frac[curr->val1] += solution[x>(*curr)];
    if (frac[curr->val1] > 0.5) {
        time[curr->val1] = curr->val2;
        //
        // Step past remaining jobs
        //
        curr_job = curr->val1;
        while (curr->val1 == curr_job)
            curr++;
    }
    else
        curr++;
}
```

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## Alternative Iteration

Computing midpoints

```
DoubleVector frac(n());
IntVector time(n());
for (int i=0; i< n(); i++){
    for (int j=EFT(i); j<LFT(i); j++){
        if (x_valid.isMember(i,j)){
            frac[j] += solution[x(i,j)];
            if (frac[j] > 0.5) {
                time[j] = j;
                continue;
            }
        }
    }
}
```

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### Building a custom indexing scheme

- Collect all the valid tuples with even indices into a list:

```
List<Tuple2<int,int>> my_index
x_const_iterator curr = x_valid().begin()
x_const_iterator end = x_valid().end()
While (curr != end) {
  if (curr->val1 % 2 == 0) && (curr->val2 % 2 == 0)
    my_index.push_back(*curr);
  curr++;
}
```

- There's no built-in iterator for this list (just an STL list)

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### Another Example: computing objective

```
DoubleVector soln(n());
IntVector index(n());
order(time,index); // sorting
soln [0] = p(index[0]);
value = p(index[0]) * w(index[0]);
for (int i=1; i< n(); i++) {
  soln[index[i]] = soln[index[i-1]] + p(index[i-1]);
  value += soln[index[i]] * w(index[i]);
}
```

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### Incumbent Heuristics

- These code fragments would be part of the `incumbentHeuristic()` method in MILPNode.
  - If you find an incumbent, this method must call `globalPtr ->updateIncumbent(solution, solution value)`
- This is a PICO (1D) solution.

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### A Customized PICO Application: A Simple Example

Command line:

```
gen_milp_app test.mod test.dat
```

Generates:

test.col	Column labels
test.row	Row labels
test.val	Values specified as PICO SYMBOLs
test.mps	The MPS file for this application
test.map	A data file that describes the sets and symbols defined by test.mod

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## A Customized PICO Application: A Simple Example

Command line:

```
gen_milp_app test.mod test.dat
```

Generates:

```
Makefile-test  generates executable
Makefile       symbolic link to Makefile-test
test_info.{h,cpp} data structures to map ampl variables
test_milp.{h,cpp} derived classes for branching, nodes, etc.
test_extras.{h,cpp} for your custom methods (never deleted)
```

Where name.{h,cpp} means name.h and name.cpp

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## Staged generation

- Stage 1 runs ampl to get row, column, val files, etc
- Stage 2 generates the C++ files with derived milp classes

You can break this up:

```
gen_milp_app -stage 1 test.mod test.dat
```

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## Syntactic Components in test.map

The AMPL/PICO interface dynamically identifies:

- Sets: a collection a n-tuples.
  - Sets of string literals: { aa, bb, cc, dd }
  - Sets of integers: { 1, 2, 3, 5, 7 }
  - Sets of n-tuples defined from other sets:  
{(aa,1), (bb,2), (cc,3), (dd,5)}
- Symbols: any row or column label

Additional sets and data values are exported from AMPL by adding to test.mod:

```
# PICO SYMBOLS: <symbol1> <symbol2> ...
```

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## Editing the map file

- Index sets are read from ampl row/col files, not model files
  - Cannot connect similar index sets

```
Set0(INT);
Set1(INT);
....
Set6(INT);
TaskDone[Set0];
UnitWork[Set1]
PrecConstraint[Set2, Set3, Set4];
X[Set5, Set6];
```

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## Editing the map file

---

- Clean up to reflect our knowledge

```
Jobs(INT);
Time(INT);
TaskDone[Jobs];
UnitWork[Time]
PrecConstraint[Jobs, Jobs, Time];
X[Jobs, Time];
```

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## Protected Symbolic Access

---

Problem: the AMPL symbols may conflict with symbols in a the PICO classes.

Solution:

- Access AMPL symbols *only* through the `info()` method  
Use `info→p(j)` instead of `p(j)`
- Enable protection by
  - Compiling with the `-DPROTECTED_APPINFO` flag
  - Running `gen_milp_app` with the `-protected-vars` flag

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## Using the Customized PICO Classes

---

By default `gen_milp_app` generates a generic `main()` function  
(parallel if you have configured with `mpi`, serial otherwise)

The following commands generate, compile and execute the PICO classes

```
gen_milp_app test.mod test.dat
make all
test_milp test.mps
```

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## Overriding Parameters

---

All parameters have default values.

```
test_milp --help
```

Should output information about parameters you can override

```
test_milp --debug=50 --useSets --preprocessIP=0 test.mps
```

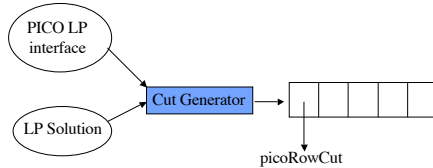
If you don't specify a value during the override, PICO will assume 1

- Fine for binary parameters
- Parameters can have any type (probably effectively any built-in type)

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## Generating Cutting Planes



- Conforms to Open Source Initiative (OSI) conventions
- Uses (and will contribute to) Cut Generation Library (CGL)
- Cuts are sparse (you have to build them)

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## ACRO

- acro-pico gives most of the pieces you need: pico, utilib, soplex, COIN
- Configuration facility
- For this workshop, you can check out from a local cvs repository
- Example configuration

Configure --site=lafayette --without-mpi --with-clp --without-cplex --with-debugging

Sets up makfiles that link properly to local MPI libraries, cplex, etc

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## Comparison with ABACUS

A Branch And Cut System

Variables and Constraints are C++ classes.

Variables must implement

```
double coeff(CONSTRAINT *c);
```

Constraints must implement

```
double coeff(VARIABLE *v);
```

Flexible, reasonably quick development, but

- User manages mapping of individual variables to linear ordering
- Cutting planes are accessed densely
- Have to compute the coefficients of all constraints, not just cuts

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